

EFFECTIVE HIRONAKA RESOLUTION AND ITS COMPLEXITY*

EDWARD BIERSTONE[†], DIMA GRIGORIEV[‡], PIERRE MILMAN[§], AND
JAROSLAW WŁODARCZYK[¶]

Dedicated to Professor Heisuke Hironaka on the occasion of his 80th birthday

Abstract. Building upon works of Hironaka, Bierstone-Milman, Villamayor and Włodarczyk. We give an *a priori* estimate for the complexity of the simplified Hironaka algorithm.

Key words. Resolution of singularities, complexity, primitive recursive functions, Grzegorzczk classes.

AMS subject classifications. 14E15, 32S45, 68Q15.

0. Introduction. In this paper we discuss the complexity of the Hironaka theorem on resolution of singularities of a marked ideal. Recall that approach to the problem of embedded resolution was originated by Hironaka (see [32]) and later developed and simplified by Bierstone-Milman (see [7], [8], [9]), Villamayor (see [49], [50]), and Włodarczyk ([53]) and others. In particular, we also use some elements from the recent development by Kollár ([38]).

It seems easier to estimate the complexity of the resolution algorithm from the recursive descriptions in Włodarczyk [53] or Bierstone-Milman [12] than from the earlier iterative versions. The algorithms in [53] and [12] (or [10]) lead to identical blowing-up sequences; whether one proof is preferable to the other is partly a matter of taste. In this article, we estimate the complexity of the “weak-strong desingularization” algorithm (see Section 1) using the construction of [53], though [12] could also be used (see Remark in this section below). In a subsequent paper, we plan to use [12] to give a comparable complexity estimate for the algorithm of “strong desingularization” (where the centres of blowing up are smooth subvarieties of the successive strict transforms).

The basic question which arises is in what terms to estimate the complexity *a priori*? We recall (see, e. g., [52], [27]) that the complexity is usually measured as a function of the bit-size of the input. In particular, in this paper we study varieties and ideals which are represented by families of polynomials with integer coefficients, and the vector of all these coefficients (for an initial variety and an ideal) is treated as an input. Hironaka’s algorithm consists of many steps of elementary calculations, but they are organized in several (nested) recursions where the resolution of an object (a variety or a marked ideal, see below) is reduced to resolutions of suitable objects

*Received April 23, 2010; accepted for publication May 20, 2011.

[†]Fields Institute, 222 College Street, Toronto, Ontario, Canada M5T 3J1; and Department of Mathematics, University of Toronto, Toronto, Ontario, Canada M5S 2E4 (bierston@fields.utoronto.ca). The author’s research was partially supported by NSERC Discovery Grant OGP 0009070.

[‡]Laboratoire des Mathématiques, Université de Lille I, 59655, Villeneuve d’Ascq, France (dmitry.grigoryev@math.univ-lille1.fr).

[§]Department of Mathematics, University of Toronto, Toronto, Ontario, Canada M5S 2E4 (milman@math.toronto.edu). The author’s research was partially supported by NSERC Discovery Grant OGP 008949.

[¶]Department of Mathematics, Purdue University, West Lafayette, IN-47907, USA (wloadar@math.purdue.edu). The author was supported in part by NSF grant DMS-0100598 and Polish grant NCN 201 608640.

with smaller values of appropriate parameters (like dimension or multiplicity). It is instructive to represent the Hironaka algorithm as a tree, to each node a of which corresponds a marked ideal. The marked ideals which correspond to child nodes of a have either smaller multiplicity of an ideal or smaller dimension of a variety. An initial marked ideal corresponds to the root of the tree. The depth of the tree is bounded by $2 \cdot m$ where m denotes the dimension of the initial variety, while the number of the nested recursions does not exceed $m + 3$. It appears that just the number of nested recursions is the overwhelming contribution to the complexity of the Hironaka's algorithm.

That is why as a relevant language for expressing a complexity bound we have chosen the Grzegorzczak class \mathcal{E}^l , $l \geq 0$ [28], [52] which consists of (integer) functions whose construction requires l nested primitive recursions. The classes \mathcal{E}^l , $l \geq 0$ provide a hierarchy of the set of all primitive-recursive functions $\cup_{l < \infty} \mathcal{E}^l$. In particular, \mathcal{E}^2 contains all (integer) polynomials and \mathcal{E}^3 contains all finite compositions of the exponential function.

As an illustration of complexity bounds from small Grzegorzczak classes, we give examples of a few algebraic-geometrical computational problems: Polynomial factoring [27], with polynomial complexity (so in \mathcal{E}^2); finding irreducible components of a variety, with the exponential complexity (so in \mathcal{E}^3) [27]; and constructing a Groebner basis of an ideal, with double-exponential complexity (so also in \mathcal{E}^3) [25, 42].

The principal complexity result of this paper (Theorem 6.4.2) states that the complexity of resolution of an ideal on an m -dimensional variety is bounded by a function from class \mathcal{E}^{m+3} . We mention also that the complexity of Hilbert's Idealbasissatz for polynomial ideals in n variables (much simpler from a purely mathematical point of view) belongs to class \mathcal{E}^{n+1} (cf. [44], [46], where the latter was formulated in different languages), and, moreover, the number $n + 1$ is sharp. This shows that these two quite different algorithmic problems have a common feature in recursion on the dimension which mainly determines their complexities.

REMARK. The main differences between the proofs in [12] and [53] come from the notions of derivative ideal that are used ([12] uses only derivatives that preserve the ideal of the exceptional divisor) and from passage to a "homogenized ideal" in [53] (see §2.8). The latter has the advantage that any two maximal contact hypersurfaces for the homogenized ideal are related by an automorphism, while [12] provides a stronger version of functoriality that is needed for strong desingularization. Since [12] does not involve homogenization, certain complexity estimates can be improved (see Remark after Corollary 5.0.10), although the overall Grzegorzczak complexity class \mathcal{E}^{m+3} is unchanged.

We mention that in [48, 47, 39] polynomial complexity algorithms for resolution of a curve are presented. Also in [13], a complexity estimated in \mathcal{E}^3 for an algorithm for resolution of singularities in the non-exceptional monomial case is presented.

In Section 1 below, we formulate the results on canonical principalization of a sheaf of ideals and on embedded desingularization. In Section 2 we give definitions of basic notions like marked ideals, hypersurfaces of maximal contact and coefficient ideals, and we formulate their properties (one can find proofs in [53]). In Section 3 we describe the resolution algorithm. In Section 4, we provide bounds on the degrees and on the number of polynomials which describe a single blow-up. In Section 5 we give some auxiliary bounds — on the multiplicity of an ideal in terms of degrees of describing polynomials, on the degree of a hypersurface of maximal contact, and on the number of generators of the coefficient ideal and their degrees. Finally in Section

6 we estimate the complexity of the resolution algorithm in terms of Grzegorzczuk's classes (their definition is also provided in Section 6).

Acknowledgements. We heartily thank the Max Planck Institut fuer Mathematik, Bonn for its warm hospitality.

1. Formulation of the Hironaka resolution theorems. All algebraic varieties in this paper are defined over a ground field of characteristic zero. The assumption of characteristic zero is only needed for the local existence of a hypersurface of maximal contact (Lemma 2.6.4).

We give proofs of the following Hironaka Theorems (see [32]):

1. Canonical Principalization.

THEOREM 1.0.1. *Let \mathcal{I} be a sheaf of ideals on a smooth algebraic variety X . There exists a principalization of \mathcal{I} ; that is, a sequence*

$$X = X_0 \xleftarrow{\sigma_1} X_1 \xleftarrow{\sigma_2} X_2 \leftarrow \dots \leftarrow X_i \leftarrow \dots \leftarrow X_r = \tilde{X}$$

of blow-ups $\sigma_i : X_{i-1} \leftarrow X_i$ of smooth centers $C_{i-1} \subset X_{i-1}$, such that:

- (a) *The exceptional divisor E_i of the induced morphism $\sigma^i = \sigma_1 \circ \dots \circ \sigma_i : X_i \rightarrow X$ has only simple normal crossings and C_i has simple normal crossings with E_i .*
- (b) *The total transform $\sigma^{r*}(\mathcal{I})$ is the ideal of a simple normal crossing divisor \tilde{E} which is a natural combination of the irreducible components of the divisor E_r .*

Moreover, the morphism $(\tilde{X}, \tilde{\mathcal{I}}) \rightarrow (X, \mathcal{I})$ defined by the above principalization commutes with smooth morphisms and embeddings of ambient varieties. It is equivariant with respect to any group action, not necessarily preserving the ground field K .

2. Weak-Strong Hironaka Embedded Desingularization.

THEOREM 1.0.2. *Let Y be a subvariety of a smooth variety X over a field of characteristic zero. There exists a sequence*

$$X_0 = X \xleftarrow{\sigma_1} X_1 \xleftarrow{\sigma_2} X_2 \leftarrow \dots \leftarrow X_i \leftarrow \dots \leftarrow X_r = \tilde{X}$$

of blow-ups $\sigma_i : X_{i-1} \leftarrow X_i$ of smooth centers $C_{i-1} \subset X_{i-1}$, such that:

- (a) *The exceptional divisor E_i of the induced morphism $\sigma^i = \sigma_1 \circ \dots \circ \sigma_i : X_i \rightarrow X$ has only simple normal crossings and C_i has simple normal crossings with E_i .*
- (b) *Let $Y_i \subset X_i$ be the strict transform of Y . All centers C_i are disjoint from the set $\text{Reg}(Y) \subset Y_i$ of points where Y (not Y_i) is smooth (and are not necessarily contained in Y_i).*
- (c) *The strict transform $\tilde{Y} := Y_r$ of Y is smooth and has only simple normal crossings with the exceptional divisor E_r .*
- (d) *The morphism $(X, Y) \leftarrow (\tilde{X}, \tilde{Y})$ defined by the embedded desingularization commutes with smooth morphisms and embeddings of ambient varieties. It is equivariant with respect to any group action, not necessarily preserving the ground K .*

3. Canonical Resolution of Singularities.

THEOREM 1.0.3. *Let Y be an algebraic variety over a field of characteristic zero. Then there exists a canonical desingularization of Y ; that is, a smooth*

variety \tilde{Y} together with a proper birational morphism $\text{res}_Y : \tilde{Y} \rightarrow Y$ such that:

- (a) $\text{res}_Y : \tilde{Y} \rightarrow Y$ is an isomorphism over the nonsingular part of Y .
- (b) The inverse image $\text{res}_Y^{-1}(Y_{\text{sing}})$ of the singular locus of Y is a simple normal crossings divisor.
- (c) The morphism res_Y is functorial with respect to smooth morphisms. For any smooth morphism $\phi : Y' \rightarrow Y$ there is a natural lifting $\tilde{\phi} : \tilde{Y}' \rightarrow \tilde{Y}$ which is a smooth morphism.
- (d) res_Y is equivariant with respect to any group action, not necessarily preserving the ground field.

REMARK. Note that a blow-up of codimension one components is an isomorphism. However it defines a nontrivial transformation of marked ideals. In the actual desingularization process, blow-ups of this kind may occur for some marked ideals induced on subvarieties of ambient varieties. Though they define isomorphisms of those subvarieties they determine blow-ups of ambient varieties which are not isomorphisms.

REMARKS. (1) By the exceptional divisor of a blow-up $\sigma : X' \rightarrow X$ with smooth center C we mean the inverse image $E := \sigma^{-1}(C)$ of the center C . By the exceptional divisor of a composite of blow-ups σ_i with smooth centers C_{i-1} we mean the union of the strict transforms of the exceptional divisors of σ_i . This definition coincides with the standard definition of the exceptional set of points of the birational morphism in the case when $\text{codim}(C_i) \geq 2$ (as in Theorem 1.0.2). If $\text{codim}(C_{i-1}) = 1$ the blow-up of C_{i-1} is an identical isomorphism and defines a formal operation of converting a subvariety $C_{i-1} \subset X_{i-1}$ into a component of the exceptional divisor E_i on X_i . This formalism is convenient for the proofs. In particular it indicates that C_{i-1} identified via σ_i with a component of E_i has simple normal crossings with other components of E_i .

(2) In Theorem 1.0.2, we blow up centers of codimension ≥ 2 and both definitions coincide.

2. Marked ideals, coefficient ideals and hypersurfaces of maximal contact. We shall assume that the ground field is algebraically closed.

2.1. Resolution of marked ideals. For any sheaf of ideals \mathcal{I} on a smooth variety X and any point $x \in X$ we denote by

$$\text{ord}_x(\mathcal{I}) := \max\{i \mid \mathcal{I} \subset m_x^i\}$$

the order of \mathcal{I} at x . (Here m_x denotes the maximal ideal of x .)

DEFINITION 2.1.1. (Hironaka (see [32], [34]), Bierstone-Milman (see [8]), Villamayor(see [49])) A *marked ideal* (originally a *basic object* in Villamayor) is a collection (X, \mathcal{I}, E, μ) , where X is a smooth variety, \mathcal{I} is a sheaf of ideals on X , μ is a nonnegative integer and E is a totally ordered collection of divisors whose irreducible components are pairwise disjoint and all have multiplicity one. Moreover the irreducible components of divisors in E have simultaneously simple normal crossings.

DEFINITION 2.1.2. (Hironaka ([32], [34]), Bierstone-Milman (see [8]), Villamayor (see [49])) By the *support* (originally *singular locus*) of (X, \mathcal{I}, E, μ) we mean

$$\text{supp}(X, \mathcal{I}, E, \mu) := \{x \in X \mid \text{ord}_x(\mathcal{I}) \geq \mu\}.$$

REMARKS. (1) Sometimes for simplicity we will represent marked ideals (X, \mathcal{I}, E, μ) as couples (\mathcal{I}, μ) or even ideals \mathcal{I} .

(2) For any sheaf of ideals \mathcal{I} on X , we have $\text{supp}(\mathcal{I}, 1) = \text{supp}(\mathcal{O}_X/\mathcal{I})$.

(3) For any marked ideal (\mathcal{I}, μ) on X , $\text{supp}(\mathcal{I}, \mu)$ is a closed subset of X (Lemma 2.5.2).

DEFINITION 2.1.3. (Hironaka (see [32], [34]), Bierstone-Milman (see [8]), Villamayor (see [49])) By a *resolution* of (X, \mathcal{I}, E, μ) we mean a sequence of blow-ups $\sigma_i : X_i \rightarrow X_{i-1}$ of smooth centers $C_{i-1} \subset X_{i-1}$,

$$X_0 = X \xleftarrow{\sigma_1} X_1 \xleftarrow{\sigma_2} X_2 \xleftarrow{\sigma_3} \dots X_i \xleftarrow{\dots} \dots \xleftarrow{\sigma_r} X_r,$$

which defines a sequence of marked ideals $(X_i, \mathcal{I}_i, E_i, \mu)$ where

1. $C_i \subset \text{supp}(X_i, \mathcal{I}_i, E_i, \mu)$.
2. C_i has simple normal crossings with E_i .
3. $\mathcal{I}_i = \mathcal{I}(D_i)^{-\mu} \sigma_i^*(\mathcal{I}_{i-1})$, where $\mathcal{I}(D_i)$ is the ideal of the exceptional divisor D_i of σ_i .
4. $E_i = \sigma_i^c(E_{i-1}) \cup \{D_i\}$, where $\sigma_i^c(E_{i-1})$ is the set of strict transforms of divisors in E_{i-1} .
5. The order on $\sigma_i^c(E_{i-1})$ is defined by the order on E_{i-1} , while D_i is the maximal element of E_i .
6. $\text{supp}(X_r, \mathcal{I}_r, E_r, \mu) = \emptyset$.

DEFINITION 2.1.4. A sequence of morphisms which are either isomorphisms or blow-ups satisfying conditions (1)-(5) is called a *multiple blow-up*. The number of morphisms in a multiple blow-up will be called its *length*.

DEFINITION 2.1.5. An *extension* of a multiple blow-up (or a resolution) $(X_i)_{0 \leq i \leq m}$ is a sequence $(X'_j)_{0 \leq j \leq m'}$ of blow-ups and isomorphisms $X'_0 = X'_{j_0} = \dots = X'_{j_1-1} \leftarrow X'_{j_1} = \dots = X'_{j_2-1} \leftarrow \dots \leftarrow X'_{j_m} = \dots = X'_{m'}$, where $X'_{j_i} = X_i$.

REMARKS. (1) The definition of extension arises naturally when we pass to open subsets of the ambient variety X .

(2) The notion of a *multiple blow-up* is analogous to the notion of a sequence of *admissible* blow-ups considered by Hironaka, Bierstone-Milman and Villamayor.

2.2. Transforms of marked ideal and controlled transforms of functions.

In the setting of the above definition we will call

$$(\mathcal{I}_i, \mu) := \sigma_i^c(\mathcal{I}_{i-1}, \mu)$$

the *transform of the marked ideal* or *controlled transform* of (\mathcal{I}, μ) . It makes sense for a single blow-up in a multiple blow-up as well as for a multiple blow-up. Let $\sigma^i := \sigma_1 \circ \dots \circ \sigma_i : X_i \rightarrow X$ be a composition of consecutive morphisms of a multiple blow-up. Then in the above setting

$$(\mathcal{I}_i, \mu) = \sigma^{ic}(\mathcal{I}, \mu).$$

We will also denote the controlled transform $\sigma^{ic}(\mathcal{I}, \mu)$ by $(\mathcal{I}, \mu)_i$ or $[\mathcal{I}, \mu]_i$.

The controlled transform can also be defined for local sections $f \in \mathcal{I}(U)$. Let $\sigma : X \leftarrow X'$ be a blow-up of a smooth center $C \subset \text{supp}(\mathcal{I}, \mu)$ defining transformation of marked ideals $\sigma^c(\mathcal{I}, \mu) = (\mathcal{I}', \mu)$. Let $f \in \mathcal{I}(U)$ be a section of \mathcal{I} . Let $U' \subseteq \sigma^{-1}(U)$ be an open subset for which the sheaf of ideals of the exceptional divisor is generated by a function y . The function

$$g = y^{-\mu}(f \circ \sigma) \in \mathcal{I}(U')$$

is the *controlled transform* of f on U' (defined up to an invertible function). As before we extend it to any multiple blow-up.

The following lemma shows that the notion of controlled transform is well defined.

LEMMA 2.2.1. *Let $C \subset \text{supp}(\mathcal{I}, \mu)$ be a smooth center of a blow-up $\sigma : X \leftarrow X'$ and let D denote the exceptional divisor. Let \mathcal{I}_C denote the sheaf of ideals defined by C . Then*

1. $\mathcal{I} \subset \mathcal{I}_C^\mu$.
2. $\sigma^*(\mathcal{I}) \subset (\mathcal{I}_D)^\mu$.

Proof. (1) We can assume that the ambient variety X is affine. Let u_1, \dots, u_k be parameters generating \mathcal{I}_C . Suppose $f \in \mathcal{I} \setminus \mathcal{I}_C^\mu$. Then we can write $f = \sum_\alpha c_\alpha u^\alpha$, where either $|\alpha| \geq \mu$ or $|\alpha| < \mu$ and $c_\alpha \notin \mathcal{I}_C$. By the assumption there is α with $|\alpha| < \mu$ such that $c_\alpha \notin \mathcal{I}_C$. Take α with the smallest $|\alpha|$. There is a point $x \in C$ for which $c_\alpha(x) \neq 0$ and in the Taylor expansion of f at x there is a term $c_\alpha(x)u^\alpha$. Thus $\text{ord}_x(\mathcal{I}) < \mu$. This contradicts to the assumption $C \subset \text{supp}(\mathcal{I}, \mu)$.

(2) $\sigma^*(\mathcal{I}) \subset \sigma^*(\mathcal{I}_C)^\mu = (\mathcal{I}_D)^\mu$. \square

2.3. Hironaka resolution principle. Our proof is based on the following principle which can be traced back to Hironaka and was used by Villamayor in his simplification of Hironaka’s algorithm:

- (1) **(Canonical) Resolution of marked ideals (X, \mathcal{I}, E, μ)**
- \Downarrow
- (2) **(Canonical) Principalization of the sheaves \mathcal{I} on X**
- \Downarrow
- (3) **(Canonical) Weak Embedded Desingularization of subvarieties $Y \subset X$**
- \Downarrow
- (4) **(Canonical) Desingularization**

(1) \Rightarrow (2). It follows immediately from the definition that a resolution of $(X, \mathcal{I}, \emptyset, 1)$ determines a principalization of \mathcal{I} . Denote by $\sigma : X \leftarrow \tilde{X}$ the morphism defined by a resolution of $(X, \mathcal{I}, \emptyset, 1)$. The controlled transform $(\tilde{\mathcal{I}}, 1) := \sigma^c(\mathcal{I}, 1)$ has the empty support. Consequently, $V(\tilde{\mathcal{I}}) = \emptyset$, and thus $\tilde{\mathcal{I}}$ is equal to the structural sheaf $\mathcal{O}_{\tilde{X}}$. This implies that the full transform $\sigma^*(\mathcal{I})$ is principal and generated by the sheaf of ideal of a divisor whose components are the exceptional divisors. The actual process of desingularization is often achieved before $(X, \mathcal{I}, E, 1)$ has been resolved (see [53]).

(2) \Rightarrow (3). Let $Y \subset X$ be an irreducible subvariety. Assume there is a principalization of sheaves of ideals \mathcal{I}_Y subject to conditions (a) and (b) of Theorem 1.0.1. Then, in the course of the principalization of \mathcal{I}_Y , the strict transform Y_i of Y in some X_i is the center of a blow-up. At this stage Y_i is nonsingular and has simple normal crossings with the exceptional divisors.

(3) \Rightarrow (4). Every algebraic variety locally admits an embedding into an affine space. Then we can show that the existence of canonical embedded desingularization independent of the embedding defines a canonical desingularization.

For more details, see [53].

REMARK. Resolution scheme and marked ideals. Marked ideals will be understood as objects which carry vital information in the resolution scheme. There are four different types of information that can be associated with marked ideals:

1. The support $\text{supp}(\mathcal{I}, \mu)$ is the “bad locus” which shall be eliminated. The blow-ups performed should have centers inside of $\text{supp}(\mathcal{I}, \mu)$.
2. The controlled transform $\sigma^c(\mathcal{I}, \mu) = \mathcal{I}_D^{-\mu} \sigma^*(\mathcal{I}, \mu)$ is the transform of the marked ideal associated with blow-ups with centers inside $\text{supp}(\mathcal{I}, \mu)$.
3. The resolution of $\text{supp}(\mathcal{I}, \mu)$ is the sequence of blow-ups and the induced transformations of marked ideals eliminating the support of the resulting marked ideal (\mathcal{I}, μ) .
4. Canonical resolution is a unique resolution which will be assigned to a marked ideal. Once we assign to a certain class of marked ideals their canonical resolutions they become useful operations to resolve some larger class of marked ideals. In other words, the resolution of a certain marked ideal is always reduced to resolution of some “simpler” marked ideals. The notion of simplicity refers essentially to two very rough invariants : the dimension of the ambient variety, and the order of nonmonomial part.

The algorithm builds upon two different canonical reductions:

- reduction of order by resolving a so called “companion ideal” (see Step 2 in Section 3).
- reduction of dimension of the ambient variety which relies on the two fundamental concepts of hypersurface of *maximal contact*, and *coefficient ideal* (see Sections 2.6, 2.9).

2.4. Equivalence relation for marked ideals. Let us introduce the following equivalence relation for marked ideals:

DEFINITION 2.4.1. Let $(X, \mathcal{I}, E_{\mathcal{I}}, \mu_{\mathcal{I}})$ and $(X, \mathcal{J}, E_{\mathcal{J}}, \mu_{\mathcal{J}})$ be two marked ideals on a smooth variety X . Then $(X, \mathcal{I}, E_{\mathcal{I}}, \mu_{\mathcal{I}}) \simeq (X, \mathcal{J}, E_{\mathcal{J}}, \mu_{\mathcal{J}})$ if:

1. $E_{\mathcal{I}} = E_{\mathcal{J}}$ and the orders on $E_{\mathcal{I}}$ and on $E_{\mathcal{J}}$ coincide.
2. $\text{supp}(\mathcal{I}, \mu_{\mathcal{I}}) = \text{supp}(\mathcal{J}, \mu_{\mathcal{J}})$.
3. The multiple blow-ups $(X_i)_{i=0, \dots, k}$ are the same for both marked ideals, and $\text{supp}(\mathcal{I}_i, \mu_{\mathcal{I}}) = \text{supp}(\mathcal{J}_i, \mu_{\mathcal{J}})$.

EXAMPLE 2.4.2. For any $k \in \mathbf{N}$, $(\mathcal{I}, \mu) \simeq (\mathcal{I}^k, k\mu)$.

REMARK. The marked ideals considered in this paper satisfy a stronger equivalence condition: For any smooth morphism $\phi : X' \rightarrow X$, $\phi^*(\mathcal{I}, \mu) \simeq \phi^*(\mathcal{J}, \mu)$. This condition will follow and is not added in the definition.

2.5. Ideals of derivatives. Ideals of derivatives were first introduced and studied in the resolution context by Giraud. Villamayor developed and applied this language to his *basic objects*.

DEFINITION 2.5.1. (Giraud, Villamayor) Let \mathcal{I} be a coherent sheaf of ideals on a smooth variety X . By the *first derivative* (originally *extension*) $\mathcal{D}(\mathcal{I})$ of \mathcal{I} we mean the coherent sheaf of ideals generated by all functions $f \in \mathcal{I}$ together with their first

derivatives. Then the i -th derivative $\mathcal{D}^i(\mathcal{I})$ is defined to be $\mathcal{D}(\mathcal{D}^{i-1}(\mathcal{I}))$. If (\mathcal{I}, μ) is a marked ideal and $i \leq \mu$ then we define

$$\mathcal{D}^i(\mathcal{I}, \mu) := (\mathcal{D}^i(\mathcal{I}), \mu - i).$$

Recall that on a smooth variety X there is a locally free sheaf of differentials $\Omega_{X/K}$ over K generated locally by du_1, \dots, du_n for a set of local parameters u_1, \dots, u_n . The dual sheaf of derivations $\text{Der}_K(\mathcal{O}_X)$ is locally generated by the derivations $\frac{\partial}{\partial u_i}$. Immediately from the definition we observe that $\mathcal{D}(\mathcal{I})$ is a coherent sheaf defined locally by generators f_j of \mathcal{I} and all their partial derivatives $\frac{\partial f_j}{\partial u_i}$. We see by induction that $\mathcal{D}^i(\mathcal{I})$ is a coherent sheaf defined locally by the generators f_j of \mathcal{I} and their derivatives $\frac{\partial^{|\alpha|} f_j}{\partial u^\alpha}$ for all multiindices $\alpha = (\alpha_1, \dots, \alpha_n)$, where $|\alpha| := \alpha_1 + \dots + \alpha_n \leq i$.

LEMMA 2.5.2. (*Giraud, Villamayor*) For any $i \leq \mu - 1$,

$$\text{supp}(\mathcal{I}, \mu) = \text{supp}(\mathcal{D}^i(\mathcal{I}), \mu - i).$$

In particular, $\text{supp}(\mathcal{I}, \mu) = \text{supp}(\mathcal{D}^{\mu-1}(\mathcal{I}), 1) = V(\mathcal{D}^{\mu-1}(\mathcal{I}))$ is a closed set.

We write $(\mathcal{I}, \mu) \subset (\mathcal{J}, \mu)$ if $\mathcal{I} \subset \mathcal{J}$.

LEMMA 2.5.3. (*Giraud, Villamayor*) Let (\mathcal{I}, μ) be a marked ideal, $C \subset \text{supp}(\mathcal{I}, \mu)$ a smooth center, and $r \leq \mu$. Let $\sigma : X \leftarrow X'$ be a blow-up at C . Then

$$\sigma^c(\mathcal{D}^r(\mathcal{I}, \mu)) \subseteq \mathcal{D}^r(\sigma^c(\mathcal{I}, \mu)).$$

Proof. See the simple computations in [51], [53]. \square

2.6. Hypersurfaces of maximal contact. The concept of the *hypersurfaces of maximal contact* is one of the key points of this proof. It was originated by Hironaka, Abhyankhar and Giraud and developed in the papers of Bierstone-Milman and Villamayor. In our terminology, we are looking for a smooth hypersurface containing the support of a marked ideal and whose strict transforms under multiple blow-ups contain the supports of the induced marked ideals. Existence of such hypersurfaces allows a reduction of the resolution problem to codimension 1.

First we introduce marked ideals which locally admit hypersurfaces of maximal contact.

DEFINITION 2.6.1. (*Villamayor* (see [49])) We say that a marked ideal (\mathcal{I}, μ) is of *maximal order* (originally *simple basic object*) if $\max\{\text{ord}_x(\mathcal{I}) \mid x \in X\} \leq \mu$ or equivalently $\mathcal{D}^\mu(\mathcal{I}) = \mathcal{O}_X$.

LEMMA 2.6.2. (*Villamayor* (see [49])) Let (\mathcal{I}, μ) be a marked ideal of maximal order and let $C \subset \text{supp}(\mathcal{I}, \mu)$ be a smooth center. Let $\sigma : X \leftarrow X'$ be a blow-up at C . Then $\sigma^c(\mathcal{I}, \mu)$ is of maximal order.

Proof. If (\mathcal{I}, μ) is a marked ideal of maximal order then $\mathcal{D}^\mu(\mathcal{I}) = \mathcal{O}_X$. Then, by Lemma 2.5.3, $\mathcal{D}^\mu(\sigma^c(\mathcal{I}, \mu)) \supset \sigma^c(\mathcal{D}^\mu(\mathcal{I}), 0) = \mathcal{O}_X$. \square

LEMMA 2.6.3. (*Villamayor* (see [49])) If (\mathcal{I}, μ) is a marked ideal of maximal order and $0 \leq i \leq \mu$, then $\mathcal{D}^i(\mathcal{I}, \mu)$ is of maximal order.

Proof. $\mathcal{D}^{\mu-i}(\mathcal{D}^i(\mathcal{I}, \mu)) = \mathcal{D}^\mu(\mathcal{I}, \mu) = \mathcal{O}_X$. \square

LEMMA 2.6.4. (*Giraud (see [23])*) Let (\mathcal{I}, μ) be a marked ideal of maximal order. Let $\sigma : X \leftarrow X'$ be a blow-up at a smooth center $C \subset \text{supp}(\mathcal{I}, \mu)$. Let $u \in \mathcal{D}^{\mu-1}(\mathcal{I}, \mu)(U)$ be a function such that, for any $x \in V(u)$, $\text{ord}_x(u) = 1$. Then

1. $V(u)$ is smooth;
2. $\text{supp}(\mathcal{I}, \mu) \cap U \subset V(u)$.

Let $U' \subset \sigma^{-1}(U) \subset X'$ be an open set where the exceptional divisor is described by y . Let $u' := \sigma^c(u) = y^{-1}\sigma^*(u)$ be the controlled transform of u . Then

1. $u' \in \mathcal{D}^{\mu-1}(\sigma^c(\mathcal{I}|_{U'}, \mu))$;
2. $V(u')$ is smooth;
3. $\text{supp}(\mathcal{I}', \mu) \cap U' \subset V(u')$
4. $V(u')$ is the restriction of the strict transform of $V(u)$ to U' .

Proof. (1) $u' = \sigma^c(u) = u/y \in \sigma^c(\mathcal{D}^{\mu-1}(\mathcal{I})) \subset \mathcal{D}^{\mu-1}(\sigma^c(\mathcal{I}))$.

(2) Since u is one of the local parameters describing the center of the blow-up, $u' = u/y$ is a parameter; that is, a function of order one.

(3) follows from (2). \square

DEFINITION 2.6.5. We will call a function

$$u \in T(\mathcal{I})(U) := \mathcal{D}^{\mu-1}(\mathcal{I}(U))$$

of multiplicity one a *tangent direction* of (\mathcal{I}, μ) on U .

As a corollary from the above we obtain the following lemma.

LEMMA 2.6.6. (*Giraud*) Let $u \in T(\mathcal{I})(U)$ be a tangent direction of (\mathcal{I}, μ) on U . Then for any multiple blow-up (U_i) of $(\mathcal{I}|_U, \mu)$, all the supports of the induced marked ideals $\text{supp}(\mathcal{I}_i, \mu)$ are contained in the strict transforms $V(u)_i$ of $V(u)$.

REMARKS. (1) Tangent directions are functions locally defining hypersurfaces of maximal contact.

(2) The main problem leading to complexity of the proofs is that of noncanonical choice of the tangent directions. We overcome this difficulty by introducing *homogenized ideals*.

2.7. Arithmetical operations on marked ideals. In this sections all marked ideals are defined for the smooth variety X and the same set of exceptional divisors E . Define the following operations of addition and multiplication of marked ideals:

1. $(\mathcal{I}, \mu_{\mathcal{I}}) + (\mathcal{J}, \mu_{\mathcal{J}}) := (\mathcal{I}^{\mu_{\mathcal{I}}} + \mathcal{J}^{\mu_{\mathcal{I}}}, \mu_{\mathcal{I}}\mu_{\mathcal{J}})$, or, more generally,

$$(\mathcal{I}_1, \mu_1) + \dots + (\mathcal{I}_m, \mu_m) := (\mathcal{I}_1^{\mu_2 \cdots \mu_m} + \mathcal{I}_2^{\mu_1 \mu_3 \cdots \mu_m} + \dots + \mathcal{I}_m^{\mu_1 \cdots \mu_{k-1}}, \mu_1 \mu_2 \cdots \mu_m)$$

(the operation of addition is not associative).

2. $(\mathcal{I}, \mu_{\mathcal{I}}) \cdot (\mathcal{J}, \mu_{\mathcal{J}}) := (\mathcal{I} \cdot \mathcal{J}, \mu_{\mathcal{I}} + \mu_{\mathcal{J}})$.

LEMMA 2.7.1. (1) $\text{supp}((\mathcal{I}_1, \mu_1) + \dots + (\mathcal{I}_m, \mu_m)) = \text{supp}(\mathcal{I}_1, \mu_1) \cap \dots \cap \text{supp}(\mathcal{I}_m, \mu_m)$. Moreover, multiple blow-ups (X_k) of $(\mathcal{I}_1, \mu_1) + \dots + (\mathcal{I}_m, \mu_m)$ are exactly those which are simultaneous multiple blow-ups for all (\mathcal{I}_j, μ_j) , and, for any k , we have the following equality for the controlled transforms $(\mathcal{I}_j, \mu_{\mathcal{I}})_k$:

$$(\mathcal{I}_1, \mu_1)_k + \dots + (\mathcal{I}_m, \mu_m)_k = [(\mathcal{I}_1, \mu_1) + \dots + (\mathcal{I}_m, \mu_m)]_k.$$

(2) $\text{supp}(\mathcal{I}, \mu_{\mathcal{I}}) \cap \text{supp}(\mathcal{J}, \mu_{\mathcal{J}}) \supseteq \text{supp}((\mathcal{I}, \mu_{\mathcal{I}}) \cdot (\mathcal{J}, \mu_{\mathcal{J}}))$. Moreover, any simultaneous multiple blow-up X_i of both ideals $(\mathcal{I}, \mu_{\mathcal{I}})$ and $(\mathcal{J}, \mu_{\mathcal{J}})$ is a multiple blow-up

for $(\mathcal{I}, \mu_{\mathcal{I}}) \cdot (\mathcal{J}, \mu_{\mathcal{J}})$, and for the controlled transforms $(\mathcal{I}_k, \mu_{\mathcal{I}})$ and $(\mathcal{J}_k, \mu_{\mathcal{J}})$, we have the equality

$$(\mathcal{I}_k, \mu_{\mathcal{I}}) \cdot (\mathcal{J}_k, \mu_{\mathcal{J}}) = [(\mathcal{I}, \mu_{\mathcal{I}}) \cdot (\mathcal{J}, \mu_{\mathcal{J}})]_k.$$

2.8. Homogenized ideals and tangent directions. Let (\mathcal{I}, μ) be a marked ideal of maximal order. Set $T(\mathcal{I}) := \mathcal{D}^{\mu-1}\mathcal{I}$. By the *homogenized ideal* we mean

$$\mathcal{H}(\mathcal{I}, \mu) := (\mathcal{H}(\mathcal{I}), \mu) = (\mathcal{I} + \mathcal{D}\mathcal{I} \cdot T(\mathcal{I}) + \dots + \mathcal{D}^i\mathcal{I} \cdot T(\mathcal{I})^i + \dots + \mathcal{D}^{\mu-1}\mathcal{I} \cdot T(\mathcal{I})^{\mu-1}, \mu)$$

REMARK. A homogenized ideal has two important properties:

1. It is equivalent to the given ideal.
2. It “looks the same” from all possible tangent directions.

By the first property we can use the homogenized ideal to construct resolution via the Giraud Lemma 2.6.6. By the second property such a construction does not depend on the choice of tangent directions.

LEMMA 2.8.1. *Let (\mathcal{I}, μ) be a marked ideal of maximal order. Then*

1. $(\mathcal{I}, \mu) \simeq (\mathcal{H}(\mathcal{I}), \mu)$ (see Definition 2.4.1).
2. For any multiple blow-up (X_k) of (\mathcal{I}, μ) ,

$$(\mathcal{H}(\mathcal{I}), \mu)_k = (\mathcal{I}, \mu)_k + [\mathcal{D}(\mathcal{I}, \mu)]_k \cdot [(T(\mathcal{I}), 1)]_k + \dots + [\mathcal{D}^{\mu-1}(\mathcal{I}, \mu)]_k \cdot [(T(\mathcal{I}), 1)]_k^{\mu-1}.$$

Although the following Lemmas 2.8.2 and 2.8.3 are used in this paper only in the case $E = \emptyset$, we formulate them in slightly more general versions.

LEMMA 2.8.2. *Let (X, \mathcal{I}, E, μ) be a marked ideal of maximal order. Assume there exist tangent directions $u, v \in T(\mathcal{I}, \mu)_x = \mathcal{D}^{\mu-1}(\mathcal{I}, \mu)_x$ at $x \in \text{supp}(\mathcal{I}, \mu)$ which are transversal to E . Then there exists an automorphism $\hat{\phi}_{uv}$ of the completion $\hat{X}_x := \text{Spec}(\hat{\mathcal{O}}_{x, X})$ such that*

1. $\hat{\phi}_{uv}^*(\mathcal{H}\hat{\mathcal{I}})_x = (\mathcal{H}\hat{\mathcal{I}})_x$;
2. $\hat{\phi}_{uv}^*(E) = E$;
3. $\hat{\phi}_{uv}^*(u) = v$;
4. $\text{supp}(\hat{\mathcal{I}}, \mu) := V(T(\hat{\mathcal{I}}, \mu))$ is contained in the fixed point set of $\hat{\phi}$.

Proof. (0) **Construction of the automorphism $\hat{\phi}_{uv}$.** Find parameters u_2, \dots, u_n transversal to u and v such that $u = u_1, u_2, \dots, u_n$ and v, u_2, \dots, u_n form two sets of parameters at x and divisors in E are described by some parameters u_i where $i \geq 2$. Set

$$\hat{\phi}_{uv}(u_1) = v, \quad \hat{\phi}_{uv}(u_i) = u_i \quad \text{for } i > 1.$$

(1) Let $h := v - u \in T(\mathcal{I})$. For any $f \in \hat{\mathcal{I}}$,

$$\hat{\phi}_{uv}^*(f) = f(u_1+h, u_2, \dots, u_n) = f(u_1, \dots, u_n) + \frac{\partial f}{\partial u_1} \cdot h + \frac{1}{2!} \frac{\partial^2 f}{\partial u_1^2} \cdot h^2 + \dots + \frac{1}{i!} \frac{\partial^i f}{\partial u_1^i} \cdot h^i + \dots$$

The latter element belongs to

$$\hat{\mathcal{I}} + \mathcal{D}\hat{\mathcal{I}} \cdot \widehat{T(\mathcal{I})} + \dots + \mathcal{D}^i\hat{\mathcal{I}} \cdot \widehat{T(\mathcal{I})}^i + \dots + \mathcal{D}^{\mu-1}\hat{\mathcal{I}} \cdot \widehat{T(\mathcal{I})}^{\mu-1} = \mathcal{H}\hat{\mathcal{I}}.$$

Hence $\hat{\phi}_{uv}^*(\hat{\mathcal{I}}) \subset \mathcal{H}\hat{\mathcal{I}}$.

(2)(3) follow from the construction.

(4) The fixed point set of $\widehat{\phi}_{uv}^*$ is defined by $u_i = \widehat{\phi}_{uv}^*(u_i)$, $i = 1, \dots, n$; that is, by $h = 0$. But $h \in \mathcal{D}^{\mu-1}(\mathcal{I})$ is 0 on $\text{supp}(\mathcal{I}, \mu)$. \square

LEMMA 2.8.3. Glueing Lemma. *Let (X, \mathcal{I}, E, μ) be a marked ideal of maximal order for which there exist tangent directions $u, v \in T(\mathcal{I}, \mu)$ at $x \in \text{supp}(\mathcal{I}, \mu)$ which are transversal to E . Then there exist étale neighborhoods $\phi_u, \phi_v : \overline{X} \rightarrow X$ of $x = \phi_u(\overline{x}) = \phi_v(\overline{x}) \in X$, where $\overline{x} \in \overline{X}$, such that*

1. $\phi_u^*(\mathcal{H}(\mathcal{I})) = \phi_v^*(\mathcal{H}(\mathcal{I}))$;
2. $\phi_u^*(E) = \phi_v^*(E)$;
3. $\phi_u^*(u) = \phi_v^*(v)$.

Set $(\overline{X}, \overline{\mathcal{I}}, \overline{E}, \mu) := \phi_u^*(X, \mathcal{H}(\mathcal{I}), E, \mu) = \phi_v^*(X, \mathcal{H}(\mathcal{I}), E, \mu)$.

(4) For any $\overline{y} \in \text{supp}(\overline{X}, \overline{\mathcal{I}}, \overline{E}, \mu)$, $\phi_u(\overline{y}) = \phi_v(\overline{y})$.

(5) For any multiple blow-up (X_i) of $(X, \mathcal{I}, \emptyset, \mu)$, the induced multiple blow-ups $\phi_u^*(X_i)$ and $\phi_v^*(X_i)$ of $(\overline{X}, \overline{\mathcal{I}}, \overline{E}, \mu)$ are the same (defined by the same centers).

Set $(\overline{X}_i) := \phi_u^*(X_i) = \phi_v^*(X_i)$.

(6) For any $\overline{y}_i \in \text{supp}(\overline{X}_i, \overline{\mathcal{I}}_i, \overline{E}_i, \mu)$, $\phi_{u_i}(\overline{y}_i) = \phi_{v_i}(\overline{y}_i)$, where $\phi_{u_i}, \phi_{v_i} : \overline{X}_i \rightarrow X_i$ are the induced morphisms.

Proof. (0) **Construction of étale neighborhoods** $\phi_u, \phi_v : U \rightarrow X$. Let $U \subset X$ be an open subset for which there exist u_2, \dots, u_n which are transversal to u and v on U such that $u = u_1, u_2, \dots, u_n$ and v, u_2, \dots, u_n form two sets of parameters on U , and the divisors in E are described by some u_i , where $i \geq 2$. Let \mathbf{A}^n be the affine space with coordinates x_1, \dots, x_n . First construct étale morphisms $\phi_1, \phi_2 : U \rightarrow \mathbf{A}^n$ with

$$\phi_1^*(x_i) = u_i \quad \text{for all } i \quad \text{and} \quad \phi_2^*(x_1) = v, \quad \phi_2^*(x_i) = u_i \quad \text{for } i > 1.$$

Then

$$\overline{X} := U \times_{\mathbf{A}^n} U$$

is a fiber product for the morphisms ϕ_1 and ϕ_2 . The morphisms ϕ_u, ϕ_v are defined to be the natural projections $\phi_u, \phi_v : \overline{X} \rightarrow U$ such that $\phi_1 \phi_u = \phi_2 \phi_v$. Set

$$\begin{aligned} w_1 &:= \phi_u^*(u) = (\phi_1 \phi_u)^*(x_1) = (\phi_2 \phi_v)^*(x_1) = \phi_v^*(v), \\ w_i &:= \phi_u^*(u_i) = \phi_v^*(u_i), \quad \text{for } i \geq 2. \end{aligned}$$

(1), (2), (3) follow from the construction.

(4) Let $h := v - u$. By the above the morphisms ϕ_u and ϕ_v coincide on $\phi_u^{-1}(V(h)) = \phi_v^{-1}(V(h))$.

By (4), a blow-up of a center $C \subset \text{supp}(\mathcal{H}(\mathcal{I}))$ lifts to the blow-ups at the same center $\phi_u^{-1}(C) = \phi_v^{-1}(C)$. Thus (5), (6) follow (see [53] for details). \square

2.9. Coefficient ideals and Giraud Lemma. The idea of a coefficient ideal was originated by Hironaka and then developed in papers of Villamayor and Bierstone-Milman.

EXAMPLE 2.9.1. Motivating example. Assume that $u = 0$ defines locally a hypersurface of maximal contact. Consider a coordinate system $u = u_1, u_2, \dots, u_n$. Write any function $f \in (\mathcal{I}, \mu)$ as follows

$$f := c_{\mu, f} \cdot u^\mu + c_{\mu-1, f}(u_2, \dots, u_n)u^{\mu-1} + \dots + c_{0, f}((u_2, \dots, u_n))$$

Then it can be easily seen that

$$\text{ord}_x(f) \geq \mu \Leftrightarrow \text{ord}_x(c_{\mu-i,f}) \geq \mu - i \quad \text{for all } i = 1, \dots, \mu.$$

In other words,

$$\text{supp}(\mathcal{I}, \mu) = \text{supp}(\text{Coeff}_{V(u)}(\mathcal{I}, \mu)),$$

where

$$\text{Coeff}_{V(u)}(\mathcal{I}, \mu) := \{(c_{\mu-i,f|V(u)})_{f \in \mathcal{I}, \mu-i} \mid i = 1, \dots, \mu\}.$$

Here $\text{Coeff}_{V(u)}(\mathcal{I}, \mu)$ can be considered as a very first definition of coefficient ideal. It allows one to reduce resolution of (\mathcal{I}, μ) to a resolution of $\text{Coeff}_{V(u)}(\mathcal{I}, \mu)$ “living” on a hypersurface of maximal contact. One of the problems here is that this definition depends on a choice of coordinates. That is why we replace it with the definition below.

It is important to observe that the controlled transformed preserves the form above:

$$\sigma^c(f, \mu) := c_{\mu,f} \cdot u'^\mu + c'_{\mu-1,f}(u'_2, \dots, u'_n)u'^{\mu-1} + \dots + c'_{0,f}(u'_2, \dots, u'_n),$$

where $u' = \sigma^c(u, 1)$, $c'_{\mu-i,f} = \sigma^c(c_{\mu-i,f}, \mu - i)$. In other words

$$\text{supp}(\sigma^c(\mathcal{I}, \mu)) = \text{supp}(\sigma^c(\text{Coeff}_{V(u)}(\mathcal{I}, \mu))).$$

The following definition modifies and generalizes the definition of Villamayor.

DEFINITION 2.9.2. Let (\mathcal{I}, μ) be a marked ideal of maximal order. By the *coefficient ideal* we mean

$$\mathcal{C}(\mathcal{I}, \mu) = \sum_{i=1}^{\mu} (\mathcal{D}^i \mathcal{I}, \mu - i).$$

REMARK. The coefficient ideal $\mathcal{C}(\mathcal{I})$ has two important properties:

1. $\mathcal{C}(\mathcal{I})$ is equivalent to \mathcal{I} .
2. The intersection of the support of (\mathcal{I}, μ) with any smooth subvariety S is the support of the restriction of $\mathcal{C}(\mathcal{I})$ to S :

$$\text{supp}(\mathcal{I}) \cap S = \text{supp}(\mathcal{C}(\mathcal{I})|_S).$$

Moreover this condition is persistent under relevant multiple blow-ups. These properties allow one to control and modify the part of support of (\mathcal{I}, μ) contained in S by applying multiple blow-ups of $\mathcal{C}(\mathcal{I})|_S$.

LEMMA 2.9.3. $\mathcal{C}(\mathcal{I}, \mu) \simeq (\mathcal{I}, \mu)$.

Proof. By Lemma 2.7.1, multiple blow-ups of $\mathcal{C}(\mathcal{I}, \mu)$ are simultaneous multiple blow-ups of $\mathcal{D}^i(\mathcal{I}, \mu)$ for $0 \leq i \leq \mu - 1$. By Lemma 2.5.3, multiple blow-ups of (\mathcal{I}, μ) define multiple blow-ups of all $\mathcal{D}^i(\mathcal{I}, \mu)$. Thus multiple blow-ups of (\mathcal{I}, μ) and $\mathcal{C}(\mathcal{I}, \mu)$ are the same and $\text{supp}(\mathcal{C}(\mathcal{I}, \mu))_k = \bigcap \text{supp}(\mathcal{D}^i \mathcal{I}, \mu - i)_k = \text{supp}(\mathcal{I}_k, \mu)$. \square

LEMMA 2.9.4. *Let (X, \mathcal{I}, E, μ) be a marked ideal of maximal order. Assume that S has only simple normal crossings with E . Then*

$$\text{supp}(\mathcal{I}, \mu) \cap S = \text{supp}(\mathcal{C}(\mathcal{I}, \mu)|_S).$$

Moreover let (X_i) be a multiple blow-up with centers C_i contained in the strict transforms $S_i \subset X_i$ of S . Then:

1. The restrictions $\sigma_{i|S_i} : S_i \rightarrow S_{i-1}$ of the morphisms $\sigma_i : X_i \rightarrow X_{i-1}$ define a multiple blow-up (S_i) of $\mathcal{C}(\mathcal{I}, \mu)|_S$.
2. $\text{supp}(\mathcal{I}, \mu) \cap S_i = \text{supp}[\mathcal{C}(\mathcal{I}, \mu)|_S]_i$.
3. Every multiple blow-up (S_i) of $\mathcal{C}(\mathcal{I}, \mu)|_S$ defines a multiple blow-up (X_i) of (\mathcal{I}, μ) with centers C_i contained in the strict transforms $S_i \subset X_i$ of $S \subset X$.

Proof. By Lemma 2.9.3, $\text{supp}(\mathcal{I}, \mu) \cap S = \text{supp}(\mathcal{C}(\mathcal{I}, \mu)) \cap S \subseteq \text{supp}(\mathcal{C}(\mathcal{I}, \mu)|_S)$. Let $x_1, \dots, x_k, y_1, \dots, y_{n-k}$ be local parameters at x such that $\{x_1 = 0, \dots, x_k = 0\}$ describes S . Then any function $f \in \mathcal{I}$ can be written as

$$f = \sum c_{\alpha f}(y) x^\alpha,$$

where $c_{\alpha f}(y)$ are formal power series in y_i .

Now $x \in \text{supp}(\mathcal{I}, \mu) \cap S$ iff $\text{ord}_x(c_\alpha) \geq \mu - |\alpha|$ for all $f \in \mathcal{I}$ and $|\alpha| \leq \mu$. Note that

$$c_{\alpha f|S} = \left(\frac{1}{\alpha!} \frac{\partial^{|\alpha|}(f)}{\partial x^\alpha} \right) \Big|_S \in \mathcal{D}^{|\alpha|}(\mathcal{I})|_S$$

and consequently $\text{supp}(\mathcal{I}, \mu) \cap S = \bigcap_{f \in \mathcal{I}, |\alpha| \leq \mu} \text{supp}(c_{\alpha f|S}, \mu - |\alpha|) \supset \text{supp}(\mathcal{C}(\mathcal{I}, \mu)|_S)$.

The above relation is preserved by multiple blow-ups of (\mathcal{I}, μ) . For details see [53]. \square

LEMMA 2.9.5. *Let $\phi : X' \rightarrow X$ be an étale morphism of smooth varieties and let $(X, \mathcal{I}, \emptyset, \mu)$ be a marked ideal. Then*

1. $\phi^*(\mathcal{D}(\mathcal{I})) = \mathcal{D}(\phi^*(\mathcal{I}))$;
2. $\phi^*(\mathcal{H}(\mathcal{I})) = \mathcal{H}(\phi^*(\mathcal{I}))$;
3. $\phi^*(\mathcal{C}(\mathcal{I})) = \mathcal{C}(\phi^*(\mathcal{I}))$.

Proof. Note that for any point $x \in X$ the completion $\widehat{\phi_x^*}$ is an isomorphism. Thus $\widehat{\phi_x^*}(\widehat{\mathcal{D}(\mathcal{I})}) = \widehat{\mathcal{D}(\phi^*(\mathcal{I}))}$ and therefore $\phi^*(\mathcal{D}(\mathcal{I})) = \mathcal{D}(\phi^*(\mathcal{I}))$. (2) and (3) follow from (1). \square

3. Resolution algorithm. The presentation of the following Hironaka resolution algorithm builds upon Bierstone-Milman’s, Villamayor’s and Włodarczyk’s algorithms which are simplifications of the original Hironaka proof. We also use Kollár’s trick allowing to completely eliminate the use of invariants.

REMARKS. (1) Note that a blow-up of codimension one components is an isomorphism. However it defines a nontrivial transformation of marked ideals. The inverse image of the center is still called the exceptional divisor.

(2) In the actual desingularization process, this kind of blow-up may occur for some marked ideals induced on subvarieties of ambient varieties. Though they define isomorphisms of these subvarieties, they determine blow-ups of ambient varieties which are not isomorphisms.

(3) The blow-up of the center C which coincides with the whole variety X is an empty set. The main feature which characterizes the notion of blow-up is the following “restriction property”:

If X is a smooth variety containing a smooth subvariety $Y \subset X$, which contains the center $C \subset Y$ then the blow-up $\sigma_{C,Y} : \tilde{Y} \rightarrow Y$ at C coincides with the strict transform of Y under the blow-up $\sigma_{C,X} : \tilde{X} \rightarrow X$, i.e.,

$$\tilde{Y} \simeq \overline{\sigma_{C,X}^{-1}(Y \setminus C)}.$$

Inductive setup. Let (X, \mathcal{I}, E, μ) denote an arbitrary marked ideal. We will present an algorithm which establishes the following assertion, by induction on the dimension of X .

THEOREM 3.0.6. *There is an associated resolution $(X_i)_{0 \leq i \leq m_X}$, called canonical, satisfying the following conditions:*

1. For any surjective locally étale morphism $\phi : X' \rightarrow X$ the induced sequence $(X'_i) = \phi^*(X_i)$ is the canonical resolution of $(X', \mathcal{I}', E', \mu) := \phi^*(X, \mathcal{I}, E, \mu)$.
2. For any locally analytic isomorphism $\phi : M' \rightarrow M$ the induced sequence $(X'_i) = \phi^*(X_i)$ is an extension of the canonical resolution of $(X', \mathcal{I}', E', \mu) := \phi^*(X, \mathcal{I}, E, \mu)$.

Proof. If $\mathcal{I} = 0$ and $\mu > 0$ then $\text{supp}(X, \mathcal{I}, \mu) = X$, and the blow-up of X is the empty set and thus it defines a unique resolution. Assume that $\mathcal{I} \neq 0$.

We will use induction on the dimension of X . If X is 0-dimensional, $\mathcal{I} \neq 0$ and $\mu > 0$ then $\text{supp}(X, \mathcal{I}, \mu) = \emptyset$ and all resolutions are trivial.

Step 1. Resolving a marked ideal (X, \mathcal{I}, E, μ) of maximal order. Before performing the resolution algorithm for the marked ideal (\mathcal{I}, μ) of maximal order in Step 1 we will replace it with the equivalent homogenized ideal $\mathcal{C}(\mathcal{H}(\mathcal{I}, \mu))$. Resolving the ideal $\mathcal{C}(\mathcal{H}(\mathcal{I}, \mu))$ defines a resolution of (\mathcal{I}, μ) at this step. To simplify notation we shall denote $\mathcal{C}(\mathcal{H}(\mathcal{I}, \mu))$ by $\mathcal{J} = (\mathcal{J}, \mu(\mathcal{J}))$.

Step 1a. Reduction to the nonboundary case. Moving $\text{supp } \mathcal{J}$ and H_α^s apart . For any multiple blow-up (X_i) of $(X, \mathcal{J}, E, \mu(\mathcal{J}))$, we will denote (for simplicity) the strict transform of E on any X_i also by E .

For any $x \in X_i$, let $s(x)$ denote the number of divisors in E through x and set

$$s_i = \max\{s(x) \mid x \in \text{supp } \mathcal{J}_i\}.$$

Let $s = s_0$. By assumption the intersections of any $s > s_0$ components of the exceptional divisors are disjoint from $\text{supp } \mathcal{J}$. Each intersection of divisors in E is locally defined by the intersection of some irreducible components of these divisors. Find all intersections H_α^s , $\alpha \in A$, of s irreducible components of divisors E such that $\text{supp } \mathcal{J} \cap H_\alpha^s \neq \emptyset$. By the maximality of s , the supports $\text{supp } \mathcal{J}_{|H_\alpha^s} \subset H_\alpha^s$ are disjoint from $H_{\alpha'}^s$, where $\alpha' \neq \alpha$.

Set

$$H^s := \bigcup_{\alpha} H_\alpha^s, \quad U^s := X \setminus H^{s+1}, \quad \underline{H}^s := H^s \setminus H^{s+1}.$$

Then $\underline{H}^s \subset U_s$ is a smooth closed subset U_s . Moreover $\underline{H}^s \cap \text{supp}(\mathcal{J}) = H^s \cap \text{supp}(\mathcal{J})$ is closed.

Construct the canonical resolution of $\mathcal{J}_{|\underline{H}^s}$. By Lemma 2.9.4, it defines a multiple blow-up of $(\mathcal{J}, \mu(\mathcal{J}))$ such that

$$\text{supp } \mathcal{J}_{j_1} \cap H_{j_1}^s = \emptyset.$$

In particular the number of the strict transforms of E passing through a single point of the support drops $s_{j_1} < s$. Now we put $s = s_{j_1}$ and repeat the procedure. We continue the above process until $s_{j_k} = s_r = 0$. Then $(X_j)_{0 \leq j \leq r}$ is a multiple blow-up of $(X, \mathcal{J}, E, \mu(\mathcal{J}))$ such that $\text{supp } \mathcal{J}_r$ does not intersect any divisor in E .

Therefore $(X_j)_{0 \leq j \leq r}$ and further longer multiple blow-ups $(X_j)_{0 \leq j \leq m}$ for any $m \geq r$ can be considered as multiple blow-ups of $(X, \mathcal{J}, \emptyset, \mu(\mathcal{J}))$ since starting from X_r the strict transforms of E play no further role in the resolution process since they do not intersect $\text{supp } \mathcal{J}_j$ for $j \geq r$. We reduce the situation to the “nonboundary case” $E = \emptyset$.

Step 1b. Nonboundary case. Let $(X_j)_{0 \leq j \leq r}$ be the multiple blow-up of $(X, \mathcal{J}, \emptyset, \mu(\mathcal{J}))$ defined in Step 1a.

For any $x \in \text{supp } \mathcal{J} \subset X$, find a tangent direction $u_\alpha \in \mathcal{D}^{\mu(\mathcal{J})-1}(\mathcal{J})$ on some neighborhood U_α of x . Then $V(u_\alpha) \subset U_\alpha$ is a hypersurface of maximal contact. By the quasicompactness of X , we can assume that the covering defined by U_α is finite. Let $U_{i\alpha} \subset X_i$ be the inverse image of U_α and let $H_{i\alpha} := V(u_\alpha)_i \subset U_{i\alpha}$ denote the strict transform of $H_\alpha := V(u_\alpha)$.

Set (see also [38])

$$\tilde{X} := \coprod U_\alpha \quad \tilde{H} := \coprod H_\alpha \subseteq \tilde{X}.$$

The closed embeddings $H_\alpha \subseteq U_\alpha$ define the closed embedding $\tilde{H} \subset \tilde{X}$ of a hypersurface of maximal contact \tilde{H} .

Consider the surjective étale morphism

$$\phi_U : \tilde{X} := \coprod U_\alpha \rightarrow X.$$

Denote by $\tilde{\mathcal{J}}$ the pull back of the ideal sheaf \mathcal{J} via ϕ_U . The multiple blow-up $(X_i)_{0 \leq i \leq r}$ of \mathcal{J} defines a multiple blow-up $(\tilde{X}_{0 \leq i \leq r})$ of $\tilde{\mathcal{J}}$ and a multiple blow-up $(\tilde{H}_i)_{0 \leq i \leq r}$ of $\tilde{\mathcal{J}}_H$.

Construct the canonical resolution of $(\tilde{H}_i)_{r \leq i \leq m}$ of the marked ideal $\tilde{\mathcal{J}}_{r|\tilde{H}_r}$ on \tilde{H}_r . It defines, by Lemma 2.9.3, a resolution $(\tilde{X}_{r \leq i \leq m})$ of $\tilde{\mathcal{J}}_r$ and thus also a resolution $(\tilde{X}_i)_{0 \leq i \leq m}$ of $(\tilde{X}, \tilde{\mathcal{J}}, \emptyset, \mu(\mathcal{J}))$. Moreover both resolutions are related by the property

$$\text{supp}(\tilde{\mathcal{J}}_i) = \text{supp}(\tilde{\mathcal{J}}_{i|\tilde{H}_i}).$$

Consider a (possible) lifting of ϕ_U :

$$\phi_{iU} : \tilde{X}_i := \coprod U_{i\alpha} \rightarrow X_i,$$

which is a surjective locally étale morphism. The lifting is constructed for $0 \leq i \leq r$.

For $r \leq i \leq m$ the resolution \tilde{X}_i is induced by the canonical resolution $(\tilde{H}_i)_{r \leq i \leq m}$ of $\tilde{\mathcal{J}}_{r|\tilde{H}_r}$.

We show that the resolution $(\tilde{X}_i)_{r \leq i \leq m}$ descends to the resolution $(X_i)_{r \leq i \leq m}$.

Let $\tilde{C}_{j_0} = \coprod C_{j_0\alpha}$ be the center of the blow-up $\tilde{\sigma}_{j_0} : \tilde{X}_{j_0+1} \rightarrow \tilde{X}_{j_0}$. The closed subset $C_{j_0\alpha} \subset U_{j_0\alpha}$ defines the center of an extension of the canonical resolution $(H_{j\alpha})_{r \leq j \leq m}$.

If $C_{j_0\alpha} \cap U_{j_0\beta} \neq \emptyset$, then, by the canonicity and condition (2) of the inductive assumption, the subset $C_{j_0\alpha\beta} := C_{j_0\alpha} \cap U_{j_0\beta}$ defines the center of an extension of of

the canonical resolution $H_{j\alpha\beta} := ((H_{j\alpha} \cap U_{j\beta}))_{r \leq j \leq m}$. On the other hand $C_{j_0\beta\alpha} := C_{j_0\beta} \cap U_{j\alpha}$ defines the center of an extension of the canonical resolution $((H_{j\beta\alpha} := H_{j\beta} \cap U_{j\alpha}))_{r \leq j \leq m}$.

By the Glueing Lemma 2.8.3 for the tangent directions u_α and u_β , there exist étale neighborhoods $\phi_{u_\alpha}, \phi_{u_\beta} : \overline{U}_{\alpha\beta} \rightarrow U_{\alpha\beta} := U_\alpha \cap U_\beta$ of $x = \phi_u(\overline{x}) = \phi_v(\overline{x}) \in X$, where $\overline{x} \in \overline{X}$, such that

1. $\phi_{u_\alpha}^*(\mathcal{J}) = \phi_{u_\beta}^*(\mathcal{J})$;
2. $\phi_{u_\alpha}^*(E) = \phi_{u_\beta}^*(E)$;
3. $\phi_{u_\alpha}^{-1}(H_{j\alpha\beta}) = \phi_{u_\beta}^{-1}(H_{j\beta\alpha})$;
4. $\phi_{u_\alpha}(\overline{x}) = \phi_{u_\beta}(\overline{x})$ for $\overline{x} \in \text{supp}(\phi_{u_\alpha}^*(\mathcal{J}))$.

Moreover, all the properties lift to the relevant étale morphisms $\phi_{u_{\alpha i}}, \phi_{u_{\beta i}} : \overline{U}_{\alpha\beta i} \rightarrow U_{\alpha\beta i}$. Consequently, by canonicity, $\phi_{u_\alpha j_0}^{-1}(C_{j_0\alpha\beta})$ and $\phi_{u_\beta j_0}^{-1}(C_{j_0\beta\alpha})$ both define the next center of the extension of the canonical resolution $\phi_{u_\alpha}^{-1}(H_{j_0\alpha\beta}) = \phi_{u_\beta}^{-1}(H_{j_0\beta\alpha})$ of $\phi_{u_\alpha j_0}^*(\mathcal{J}|_{H_{\alpha\beta}}) = \phi_{u_\beta}^*(\mathcal{J}|_{H_{\beta\alpha}})$. Thus

$$\phi_{u_\alpha}^{-1}(C_{j_0\alpha\beta}) = \phi_{u_\beta}^{-1}(C_{j_0\beta\alpha}),$$

and finally, by property (4),

$$C_{j_0\alpha\beta} = C_{j_0\beta\alpha}.$$

Consequently \widetilde{C}_{j_0} descends to a smooth closed center $C_{j_0} = \bigcup C_{j_0\alpha} \subset X_{j_0}$ and the resolution $(\widetilde{X}_i)_{r \leq i \leq m}$ descends to a resolution $(X_i)_{r \leq i \leq m}$.

Step 2. Resolving marked ideals (X, \mathcal{I}, E, μ) . For any marked ideal (X, \mathcal{I}, E, μ) write

$$I = \mathcal{M}(\mathcal{I})\mathcal{N}(\mathcal{I}),$$

where $\mathcal{M}(\mathcal{I})$ is the *monomial part* of \mathcal{I} , that is, the product of the principal ideals defining the irreducible components of the divisors in E , and $\mathcal{N}(\mathcal{I})$ is the *nonmonomial part* which is not divisible by any ideal of a divisor in E . Let

$$\text{ord}_{\mathcal{N}(\mathcal{I})} := \max\{\text{ord}_x(\mathcal{N}(\mathcal{I})) \mid x \in \text{supp}(\mathcal{I}, \mu)\}.$$

DEFINITION 3.0.7. (Hironaka, Bierstone-Milman, Villamayor, Encinas-Hauser) By the *companion ideal* of (\mathcal{I}, μ) where $I = \mathcal{N}(\mathcal{I})\mathcal{M}(\mathcal{I})$ we mean the marked ideal of maximal order

$$O(\mathcal{I}, \mu) = \begin{cases} (\mathcal{N}(\mathcal{I}), \text{ord}_{\mathcal{N}(\mathcal{I})}) + (\mathcal{M}(\mathcal{I}), \mu - \text{ord}_{\mathcal{N}(\mathcal{I})}) & \text{if } \text{ord}_{\mathcal{N}(\mathcal{I})} < \mu, \\ (\mathcal{N}(\mathcal{I}), \text{ord}_{\mathcal{N}(\mathcal{I})}) & \text{if } \text{ord}_{\mathcal{N}(\mathcal{I})} \geq \mu. \end{cases}$$

In particular $O(\mathcal{I}, \mu) = (\mathcal{I}, \mu)$ for ideals (\mathcal{I}, μ) of maximal order.

Step 2a. Reduction to the monomial case by using companion ideals.

By Step 1 we can resolve the marked ideal of maximal order $O(\mathcal{I}) = O(\mathcal{I}, \mu) = (O(\mathcal{I}), \mu(O(\mathcal{I})))$. By Lemma 2.7.1, for any multiple blow-up of $O(\mathcal{I}, \mu)$,

$$\begin{aligned} \text{supp}(O(\mathcal{I}, \mu))_i &= \text{supp}[\mathcal{N}(\mathcal{I}), \text{ord}_{\mathcal{N}(\mathcal{I})}]_i \cap \text{supp}[\mathcal{M}(\mathcal{I}), \mu - \text{ord}_{\mathcal{N}(\mathcal{I})}]_i \\ &= \text{supp}[\mathcal{N}(\mathcal{I}), \text{ord}_{\mathcal{N}(\mathcal{I})}]_i \cap \text{supp}(\mathcal{I}, \mu)_i. \end{aligned}$$

Consequently, such a resolution leads to an ideal (\mathcal{I}_{r_1}, μ) such that $\text{ord}_{\mathcal{N}(\mathcal{I}_{r_1})} < \text{ord}_{\mathcal{N}(\mathcal{I})}$. Then we repeat the procedure for (\mathcal{I}_{r_1}, μ) . We find marked ideals $(\mathcal{I}_{r_0}, \mu) = (\mathcal{I}, \mu), (\mathcal{I}_{r_1}, \mu), \dots, (\mathcal{I}_{r_m}, \mu)$ such that $\text{ord}_{\mathcal{N}(\mathcal{I}_0)} > \text{ord}_{\mathcal{N}(\mathcal{I}_{r_1})} > \dots > \text{ord}_{\mathcal{N}(\mathcal{I}_{r_m})}$. The procedure terminates after a finite number of steps when we arrive at an ideal (\mathcal{I}_{r_m}, μ) with $\text{ord}_{\mathcal{N}(\mathcal{I}_{r_m})} = 0$ or with $\text{supp}(\mathcal{I}_{r_m}, \mu) = \emptyset$. In the second case we get a resolution. In the first case $\mathcal{I}_{r_m} = \mathcal{M}(\mathcal{I}_{r_m})$ is monomial.

Step 2b. Monomial case $\mathcal{I} = \mathcal{M}(\mathcal{I})$. The collection of divisors E is ordered (see Definitions 2.1.1 and 2.1.3); say $E = \{D_1, D_2, \dots\}$. Let $\text{Sub}(E)$ denote the set of all subsets of E . The ordering of E induces a natural lexicographic order on $\text{Sub}(E)$: We can associate to each $S \in \text{Sub}(E)$ the lexicographic order of the binary sequence $(\delta_1, \delta_2, \dots)$, where $\delta_i = 0$ or 1 according as $D_i \notin S$ or $D_i \in S$. (The actual formula for the order is irrelevant as long as it is canonical and linear for the divisors passing through a point.)

Let x_1, \dots, x_k define equations of the components $D_1^x, \dots, D_k^x \in E$ through $x \in \text{supp}(X, \mathcal{I}, E, \mu)$ and let \mathcal{I} be generated by a monomial $x^{(a_1, \dots, a_k)}$ at x . In particular

$$\text{ord}_x(\mathcal{I})(x) := a_1 + \dots + a_k.$$

Let $\rho(x) := \{D_{i_1}, \dots, D_{i_l}\} \in \text{Sub}(E)$ be the maximal (with respect to the order on $\text{Sub}(E)$) subset satisfying the properties

1. $a_{i_1} + \dots + a_{i_l} \geq \mu$.
2. For any $j = 1, \dots, l, a_{i_1} + \dots + \check{a}_{i_j} + \dots + a_{i_l} < \mu$.

Let $R(x)$ denote the subsets in $\text{Sub}(E)$ satisfying the properties (1) and (2). The maximal irreducible components of the $\text{supp}(\mathcal{I}, \mu)$ through x are described by the intersections $\bigcap_{D \in A} D$ where $A \in R(x)$. In particular $\text{supp}(\mathcal{I}, \mu)$ is a union of components with simple normal crossings.

The maximal locus of ρ determines at most one maximal component of $\text{supp}(\mathcal{I}, \mu)$ through each x . The invariant ρ is introduced to describe the center of the blow-up in a unique way. As we see below to resolve monomial case we can randomly pick any maximal irreducible component of $\text{supp}(\mathcal{I}, \mu)$. The algorithm is controlled by the order and there is no need to introduce additional invariants unless we would like to construct invariants which describe the center and decrease after blow-up. (see [10], [20])

After the blow-up at the maximal locus $C = \{x_{i_1} = \dots = x_{i_l} = 0\}$ of ρ , the ideal $\mathcal{I} = (x^{(a_1, \dots, a_k)})$ is equal to $\mathcal{I}' = (x'^{(a_1, \dots, a_{i_j-1}, a, a_{i_j+1}, \dots, a_k)})$ in the neighborhood corresponding to x_{i_j} , where $a = a_{i_1} + \dots + a_{i_l} - \mu < a_{i_j}$. In particular the invariant $\text{ord}_x(\mathcal{I})$ drops for all points of some maximal components of $\text{supp}(\mathcal{I}, \mu)$. Thus the maximal value of $\text{ord}_x(\mathcal{I})$ on the maximal components of $\text{supp}(\mathcal{I}, \mu)$ which were blown up is bigger than the maximal value of $\text{ord}_x(\mathcal{I})$ on the new maximal components of $\text{supp}(\mathcal{I}, \mu)$. The algorithm terminates after a finite number of steps. \square

3.1. Summary of the resolution algorithm. The resolution algorithm can be represented by the following scheme.

Step 2. Resolve (\mathcal{I}, μ) .

Step 2a. Reduce (\mathcal{I}, μ) to the monomial marked ideal $\mathcal{I} = \mathcal{M}(\mathcal{I})$.

\Downarrow

If $\mathcal{I} \neq \mathcal{M}(\mathcal{I})$, decrease the maximal order of the nonmonomial part $\mathcal{N}(\mathcal{I})$ by resolving the companion ideal $O(\mathcal{I}) = O(\mathcal{I}, \mu)$.

Step 1. Resolve the companion ideal $O(\mathcal{I})$:
Replace $O(\mathcal{I})$ with $\mathcal{J} := \mathcal{C}(\mathcal{H}(O(\mathcal{I}))) \simeq O(\mathcal{I})$. (*)

Step 1a. Move apart all strict transforms of E and $\text{supp } \mathcal{J}$.

↓

Move away all intersections H_α^s of s divisors in E (where s is the maximal number of divisors in E through points in $\text{supp } \mathcal{J} = \text{supp } O(\mathcal{I})$).

↕

For any α , resolve $\mathcal{J}|_{(\cup_\alpha H_\alpha^s)}$.

Step 1b. If the strict transforms of E do not intersect $\text{supp } \mathcal{J}$, resolve \mathcal{J} .

↕

Simultaneously resolve all $\mathcal{J}|_{V(u)}$, where $V(u)$ is a hypersurface of maximal contact. (Use the property of homogenization ([53]), and Kollár's trick ([38]).

Step 2b. Resolve the monomial marked ideal $\mathcal{I} = \mathcal{M}(\mathcal{I})$.

REMARKS. (1) (*) The ideal $O(\mathcal{I})$ is replaced with $\mathcal{H}(O(\mathcal{I}))$ to ensure that the algorithm constructed in Step 1b is independent of the choice of the tangent direction u .

We replace $\mathcal{H}(O(\mathcal{I}))$ with $\mathcal{C}(\mathcal{H}(O(\mathcal{I})))$ to ensure the equalities $\text{supp}(\mathcal{J}|_S) = \text{supp}(\mathcal{J}) \cap S$, where $S = H_\alpha^s$ in Step 1a and $S = V(u)$ in Step 1b.

(2) If $\mu = 1$ the companion ideal is equal to $O(\mathcal{I}, 1) = (\mathcal{N}(\mathcal{I}), \mu_{\mathcal{N}(\mathcal{I})})$ so the general strategy of the resolution of (\mathcal{I}, μ) is to decrease the order of the nonmonomial part and then to resolve the monomial part.

(3) In particular, in order to desingularize Y , we put $\mu = 1$ and $\mathcal{I} = \mathcal{I}_Y$, where \mathcal{I}_Y is the sheaf of the subvariety Y , and we resolve the marked ideal $(X, \mathcal{I}, \emptyset, \mu)$. The nonmonomial part $\mathcal{N}(\mathcal{I}_i)$ is nothing but the weak transform $(\sigma^i)^w(\mathcal{I})$ of \mathcal{I} .

In the next sections, we provide a complexity bound for the algorithm.

4. Complexity bounds on a blow-up. Our purpose for the rest of the paper is to estimate the complexity of the desingularization algorithm described in the previous sections.

4.1. Preliminary setup.

4.1.1. Affine marked ideals. An input of the algorithm is an *affine marked ideal*; that is, a collection of tuples

$$\mathcal{T} := (\{X_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}, E_{\alpha,\beta}, U_{\alpha,\beta}, (\mathbb{C}^{n_\alpha})_\alpha \mid \alpha \in A, \beta \in B_\alpha\}, \mu),$$

where:

1. $(\mathbb{C}^{n_\alpha})_\alpha \simeq \mathbb{C}^{n_\alpha}$.
2. $\{U_{\alpha,\beta} \mid \beta \in B_\alpha\}$ is an open cover of $(\mathbb{C}^{n_\alpha})_\alpha$.
3. $U_{\alpha,\beta} \subset (\mathbb{C}^{n_\alpha})_\alpha$ is an open subset whose complement is given by $f_{\alpha,\beta} = 0$.

4. $X_{\alpha,\beta} \subset (\mathbb{C}^{n_\alpha})_\alpha$ is a closed subset such that $X_{\alpha,\beta} \cap U_{\alpha,\beta}$ is a nonsingular m -dimensional variety (possibly reducible). Moreover there exists a set of parameters (coordinates) on $U_{\alpha,\beta}$,

$$u_{\alpha,\beta,1}, \dots, u_{\alpha,\beta,n_\alpha} \in \mathbb{C}[x_{\alpha,1}, \dots, x_{\alpha,n_\alpha}],$$

such that $u_{\alpha,\beta,i}$ is a coordinate $x_{\alpha,j}$ describing an exceptional divisor or it is transversal to the exceptional divisors (over $U_{\alpha,\beta}$), and, moreover, $\mathcal{I}_{X_{\alpha,\beta}} = (u_{\alpha,\beta,i_1}, \dots, u_{\alpha,\beta,i_k}) \subset \mathbb{C}[x_{\alpha,1}, \dots, x_{\alpha,n_\alpha}]$, for a certain subset $\{i_1, \dots, i_k\} \subset \{1, \dots, n_\alpha\}$.

5. $\mathcal{I}_{\alpha,\beta} = \langle g_{\alpha,\beta,1}, \dots, g_{\alpha,\beta,j} \rangle \subset \mathbb{C}[x_{\alpha,1}, \dots, x_{\alpha,n_\alpha}]$ is an ideal,
 6. $E_{\alpha,\beta}$ is a collection of $s \leq m$ smooth divisors in $(\mathbb{C}^{n_\alpha})_\alpha$ described by some $x_{\alpha,j} = 0$, where $j = n_\alpha - s + 1, \dots, n_\alpha$. The divisors in $E_{\alpha,\beta}$ are either transversal to $X_{\alpha,\beta}$ or they contain $X_{\alpha,\beta}$. The restrictions of the divisors in $E_{\alpha,\beta}$ that are transversal to $X_{\alpha,\beta}$ define the exceptional divisors on $X_{\alpha,\beta}$.
 7. There exist birational maps $i_{\alpha_1\beta_1, \alpha_2\beta_2} : X_{\alpha_1\beta_1} \dashrightarrow X_{\alpha_2\beta_2}$ given by

$$X_{\alpha_1\beta_1} \ni x \mapsto i_{\alpha_1\beta_1, \alpha_2\beta_2}(x) = \left(\frac{v_{\alpha_1\beta_1, \alpha_2\beta_2, 1}}{w_{\alpha_1\beta_1, \alpha_2\beta_2, 1}}, \dots, \frac{v_{\alpha_1\beta_1, \alpha_2\beta_2, n_{\alpha_2}}}{w_{\alpha_1\beta_1, \alpha_2\beta_2, n_{\alpha_2}}} \right)(x) \in X_{\alpha_2\beta_2}$$

for regular functions $v_{\alpha_1\beta_1, \alpha_2\beta_2, 1}, \dots, v_{\alpha_1\beta_1, \alpha_2\beta_2, n_{\alpha_2}}$,

$w_{\alpha_1\beta_1, \alpha_2\beta_2, 1}, \dots, w_{\alpha_1\beta_1, \alpha_2\beta_2, n_{\alpha_2}}$ in $\mathbb{C}[x_{\alpha_1,1}, \dots, x_{\alpha_1, n_{\alpha_1}}]$.

8. The birational maps $i_{\alpha\beta, \alpha', \beta'}$ determine uniquely up to an isomorphism a variety $X_{\mathcal{T}}$ in the following sense: There exist open embeddings $j_{\alpha,\beta} : X_{\alpha,\beta} \cap U_{\alpha,\beta} \hookrightarrow X_{\mathcal{T}}$ defining an open cover of $X_{\mathcal{T}}$, and satisfying

$$j_{\alpha_2\beta_2}^{-1} j_{\alpha_1\beta_1} = i_{\alpha_1\beta_1, \alpha_2\beta_2}.$$

9. $\mu \geq 0$ is an integer.
 10. $\text{supp}(\mathcal{I}_{\alpha,\beta}, \mu) \cap U_{\alpha,\beta} \cap U_{\alpha,\beta'} = \text{supp}(\mathcal{I}_{\alpha,\beta'}, \mu) \cap U_{\alpha,\beta} \cap U_{\alpha,\beta'}$.

REMARKS. (1) The objects $X_{\alpha,\beta}, E_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}$, as well as corresponding functions $g_{\alpha,\beta,i}, u_{\alpha,\beta,j}, x_{\alpha,i}$ are relevant for the algorithm after they are restricted to $U_{\alpha,\beta}$. Their behavior in the complement $(\mathbb{C}^{n_\alpha})_\alpha \setminus U_{\alpha,\beta}$ has no relevance.

(2) The operation of restricting to maximal contact leads to considering open subsets $U_{\alpha,\beta} \subset (\mathbb{C}^{n_\alpha})_\alpha$.

(3) While studying the complexity of the algorithm we will assume that the coefficients of the input polynomials belong just to \mathbb{Z} . All general considerations are given for coefficients from \mathbb{C} , and remain valid for any algebraically closed field of zero characteristic.

(4) The open embeddings $j_{\alpha,\beta}$ can be constructed from $i_{\alpha_1\beta_1, \alpha_2\beta_2}$ after performing the algorithm but we do not dwell on this.

DEFINITION 4.1.1. By the support of \mathcal{T} we mean the collection of the sets

$$\text{supp}(\mathcal{T}) := \{\text{supp}(X_{\alpha,\beta} \cap U_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}, E_{\alpha,\beta} \cap U_{\alpha,\beta} \cap X_{\alpha,\beta}, \mu)\}_{\alpha \in A, \beta \in B}$$

DEFINITION 4.1.2. Given an affine marked ideal $\mathcal{T} := (\{X_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}, E_{\alpha,\beta}, U_{\alpha,\beta}, (\mathbb{C}^{n_\alpha})_\alpha \mid \alpha \in A, \beta \in B_\alpha\}, \mu)$, we say that an affine marked ideal $\mathcal{T}' := (\{X_{\alpha',\beta'}, \mathcal{I}_{\alpha',\beta'}, E_{\alpha',\beta'}, U_{\alpha',\beta'}, (\mathbb{C}^{n_\alpha})_\alpha \mid \alpha \in A', \beta \in B'_\alpha\}, \mu)$, is defined over \mathcal{T} , provided:

1. There exist maps of index-sets $p : A' \rightarrow A$, and $p_{\alpha'} : B'_{\alpha'} \rightarrow B_{p(\alpha')}$.
2. The canonical projection on the first $\alpha = p(\alpha')$ components

$$\pi'_{\alpha} : (\mathbb{C}^{n_{\alpha'}})_{\alpha'} \rightarrow (\mathbb{C}^{n_{\alpha}})_{\alpha}$$

determine birational morphisms $\pi_{\alpha',\beta'} := \pi'_{X_{\alpha',\beta'}} : X_{\alpha',\beta'} \rightarrow X_{\alpha,\beta}$ commuting with $i_{\alpha_1\beta_1,\alpha_2,\beta_2}$ and $i_{\alpha'_1\beta'_1,\alpha'_2,\beta'_2}$.

3. There exist natural birational morphisms $X_{\mathcal{T}'} \rightarrow X_{\mathcal{T}}$ commuting with $j_{\alpha,\beta}$, and $j_{\alpha',\beta'}$.

We introduce the following functions to characterize the affine marked ideal

$$\mathcal{T} := (\{X_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}, E_{\alpha,\beta}, U_{\alpha,\beta}\}_{\alpha,\beta}, (\mathbb{C}^{n_{\alpha}})_{\alpha}, \mu) :$$

1. $m(\mathcal{T}) = \dim(X_{\mathcal{T}})$,
2. $\mu(\mathcal{T}) = \mu$,
3. $d(\mathcal{T})$ is the maximal degree of all polynomials in

$$\Psi(\mathcal{T}) := \{u_{\alpha,\beta,i}, g_{\alpha,\beta,i}, f_{\alpha,\beta}, v_{\alpha,\beta,\alpha_2,\beta_2,i}, w_{\alpha,\beta,\alpha_2,\beta_2,j}\}$$

4. $n(\mathcal{T}) = \max n_{\alpha}$,
5. $l(\mathcal{T})$ is the maximal number of all polynomials in $\Psi(\mathcal{T})$.
6. $q(\mathcal{T})$ is the number of neighborhoods $U_{\alpha,\beta}$ in \mathcal{T} (i. e the number of the indices such that $\alpha \in A, \beta \in B$).
7. $b(\mathcal{T})$ the maximum bit size of any (integer) coefficient of each of the polynomials in $\Psi(\mathcal{T})$

REMARK. The function $b(\mathcal{T})$ is used only for the estimation of the total complexity of the algorithm. In particular it has no relevance for the estimates of the number of blow-ups, the maximal embedding dimension $n(\mathcal{T})$, or the number of neighborhoods.

Algorithmically the input is represented by the coefficients of polynomials describing an affine marked ideal \mathcal{T}_0 . We assume

$$m(\mathcal{T}_0) = m, \quad d(\mathcal{T}_0) \leq d_0, \quad n(\mathcal{T}_0) \leq n_0, \quad l(\mathcal{T}_0) \leq l_0, \quad b(\mathcal{T}_0) \leq b_0.$$

Then in particular, the total bit-size of the input does not exceed $b_0 \cdot l_0 \cdot d_0^{O(n_0)}$, cf. [27].

4.1.2. Resolution of singularities. For simplicity consider an irreducible affine variety $Y \subset \mathbb{C}^n$ described by some equations. The algorithm resolves Y by the following procedure:

Step A. Find the generators of $\mathcal{I}_Y = \langle g_1, \dots, g_{\bar{j}} \rangle \subset \mathbb{C}[x_1, \dots, x_n]$ and construct the affine marked ideal

$$\mathcal{T} := (X = \mathbb{C}^n, \mathcal{I}_Y, E = \emptyset, U = \mathbb{C}^n, \mathbb{C}^n, \mu = 1)$$

Step B. Start the resolving procedure for the affine marked ideal $(\{X = \mathbb{C}^n, \mathcal{I}_Y, E = \emptyset, U = \mathbb{C}^n\}, \mathbb{C}^n, \mu = 1)$ (see below).

Step C. Pick a nonsingular point $p \in Y \subset \mathbb{C}^n$. Stop the resolution procedure when the constructed center of the following blow-up in the algorithm passes through

the inverse image of p . As an output of the resolution algorithm we get an affine marked ideal

$$\mathcal{T}' := (\{X_{\alpha',\beta'}, \mathcal{I}_{\alpha',\beta'}, E_{\alpha',\beta'}, U_{\alpha,\beta}, \mathbb{C}_{\alpha'}^n\}_{\alpha',\beta'}, 1)$$

over \mathcal{T} . In particular we have a collection of projections

$$\pi_{\alpha'} : (\mathbb{C}^{n_{\alpha'}})_{\alpha'} \rightarrow \mathbb{C}^n$$

(projection on the first n coordinates.). (Note that the restriction of π_{α} defines a birational morphism $\pi_{\alpha',\beta'} : X_{\alpha',\beta'} \rightarrow X$ which is an isomorphism in a neighborhood of $p \in X$.)

The center of the following blow-up is described on some open subcover $\{U_{\alpha',\beta''}\}_{\alpha' \in A', \beta'' \in B''_{\alpha'}}$ of $\{U_{\alpha',\beta'}\}_{\alpha' \in A', \beta' \in B'_{\alpha'}}$ by $C_{\alpha',\beta''} \cap U_{\alpha',\beta''}$, for closed subsets $C_{\alpha',\beta''} \subset (\mathbb{C}^{n_{\alpha'}})_{\alpha'}$. Consider the unique irreducible component $\tilde{C}_{\alpha',\beta''}$ of $C_{\alpha',\beta''}$ containing the inverse image of the point p . Then

$$\bar{\pi}_{\alpha',\beta''} := \pi|_{\tilde{C}_{\alpha',\beta''} \cap U_{\alpha',\beta''}} : \tilde{Y}_{\alpha',\beta''} := \tilde{C}_{\alpha',\beta''} \cap U_{\alpha',\beta''} \rightarrow Y$$

is a local resolution of Y . The resolution space \tilde{Y} is described by an open cover $\{\tilde{Y}_{\alpha',\beta''}, \bar{\pi}_{\alpha',\beta''}\}_{\alpha',\beta''}$. The sets $\tilde{Y}_{\alpha',\beta''}$ are represented as closed subsets of the open subsets $U_{\alpha',\beta''} \subset \mathbb{C}^{n_{\alpha'}}$.

4.1.3. Principalization. Given a smooth affine variety $X \subset \mathbb{C}^n$, described by equations $u_{X,i} = 0$, and an ideal $\mathcal{I} = (g_1, \dots, g_k)$ on X and on \mathbb{C}^n .

Step A. First the algorithm finds affine neighbourhoods $U_{\alpha,\beta}$ in \mathbb{C}^n each given by an inequality $f_{\alpha,\beta} \neq 0$ in which X is represented by a family of local parameters

$$u_1 = \dots = u_{n-m} = 0.$$

Moreover, it finds coordinates x_i on \mathbb{C}^n , such that (up to index permutation)

$$u_1, \dots, u_{n-m}, x_{n-m+1}, \dots, x_n$$

is a complete set of parameters.

The local parameters u_i are chosen among the input polynomials. To this end the algorithm can for each choice of $\{i_1, \dots, i_{n-m}\} \subset \{1, \dots, n\}$ pick an identically non-vanishing minor $f_{\alpha,\beta}$ of the Jacobian matrix of $\{u_i\}_i$.

We construct an affine marked ideal given by an input tuple by

$$\mathcal{T} := (\{X_{\beta} := X, \mathcal{I}_{\beta} := \mathcal{I}, E_{\beta} = \emptyset, U_{\beta}, \mathbb{C}^n\}_{\beta}, \mu = 1).$$

Step B. The algorithm resolves $\mathcal{T} = (\{X_{\beta} := X, \mathcal{I}_{\beta} := \mathcal{I}, E_{\beta} = \emptyset, U_{\beta}, \mathbb{C}^n\}_{\beta}, \mu = 1)$. As an output, we get

$$\mathcal{T}' := (\{X_{\alpha',\beta'}, \mathcal{I}_{\alpha',\beta'}, E_{\alpha',\beta'}, U_{\alpha',\beta'}, (\mathbb{C}^{n_{\alpha'}})_{\alpha'}\}_{\alpha',\beta'}, 1),$$

over \mathcal{T} .

Step C. The variety $X' := X_{\mathcal{T}'}$ is described by an open cover $\{X_{\alpha',\beta'} \cap U_{\alpha',\beta'}\}_{\alpha',\beta'}$ for closed subsets $X_{\alpha',\beta'} \subset \mathbb{C}^{n_{\alpha'}}$, and open subsets $U_{\alpha',\beta'} \subset \mathbb{C}^{n_{\alpha'}}$ (see (8) from 4.1.1). Moreover, we have a collection of birational morphisms $\pi_{\alpha',\beta'} : X_{\alpha',\beta'} \cap U_{\alpha',\beta'} \rightarrow X \subset \mathbb{C}^n$. The principal ideal on $X_{\alpha',\beta'}$ is generated by

$$(g_1 \circ \pi_{\alpha',\beta'} \dots, g_k \circ \pi_{\alpha',\beta'}) = (x_1^{a_1} \dots x_n^{a_n}).$$

4.2. Description of blow-up. Consider an affine marked ideal

$$\mathcal{T} := (\{X_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}, E_{\alpha,\beta}, U_{\alpha,\beta}, (\mathbb{C}^{n_\alpha})_\alpha \mid \alpha \in A, \beta \in B_\alpha\}, \mu)$$

corresponding to a marked ideal (X, \mathcal{I}, E, μ) . Let $C \subset X$ be a smooth center described as follows:

We assume that there is an open subcover $\{U_{\alpha,\beta'}\}_{\alpha \in A, \beta' \in B'_\alpha} \subset (\mathbb{C}^{n_\alpha})_\alpha \simeq \mathbb{C}^n$ of $U_{\alpha,\beta}$, together with a map of indices $\rho : B'_\alpha \rightarrow B_\alpha$, and a collection of closed subvarieties $C_{\alpha,\beta'} \subset (\mathbb{C}^{n_\alpha})_\alpha$ (of dimension $k_{\alpha,\beta'} \leq m$), such that

1. $\bigcup_{\rho(\beta')=\beta} U_{\alpha,\beta'} = U_{\alpha,\beta}$;
2. $C_{\alpha,\beta'} \cap U_{\alpha,\beta'} \subset \text{supp}(\mathcal{I}_{\alpha,\beta',\mu}) \cap U_{\alpha,\beta'}$;
3. $C_{\alpha,\beta'}$ is described on each $U_{\alpha,\beta'}$ by a set of local parameters

$$u_{\alpha,\beta',1}, \dots, u_{\alpha,\beta',n_\alpha-m}, u_{\alpha,\beta',n_\alpha-m+1}, \dots, u_{\alpha,\beta',n_\alpha-k_{\alpha,\beta'}} \in \mathbb{C}[x_1, \dots, x_{n_\alpha}],$$

i.e.,

$$u_{\alpha,\beta',1} = \dots = u_{\alpha,\beta',n_\alpha-m} = u_{\alpha,\beta',n_\alpha-m+1} = \dots = u_{\alpha,\beta',n_\alpha-k_{\alpha,\beta'}} = 0,$$

where $X_{\alpha,\beta}$ is described on $U_{\alpha,\beta} \supset U_{\alpha,\beta'}$ by $u_{\alpha,\beta',1} = \dots = u_{\alpha,\beta',n_\alpha-m} = 0$;

4. $u_{\alpha,\beta',1}, \dots, u_{\alpha,\beta',n_\alpha-k_{\alpha,\beta'}}$ are transversal to the exceptional divisors (over $U_{\alpha,\beta'}$), or coincide with coordinate functions describing the exceptional divisors.

Denote by

$$\mathcal{T}' := (\{X_{\alpha',\beta'}, \mathcal{I}_{\alpha',\beta'}, E_{\alpha',\beta'}, U_{\alpha',\beta'}, (\mathbb{C}^{n_{\alpha'}})_{\alpha'} \mid \alpha' \in A', \beta' \in B'_{\alpha'}\}, \mu)$$

the resulting affine marked ideal obtained from \mathcal{T} by the blow-up with the center C . Below we describe more precisely the ingredients of \mathcal{T}' .

The open cover after blow-up. The blow-up creates a new collection of ambient affine spaces $(\mathbb{C}^{n_{\alpha'}})_{\alpha'}$. Namely, we can associate with functions $u_{\alpha,\beta',i}$ on $(\mathbb{C}^{n_\alpha})_\alpha$, where $i = 1, \dots, n_\alpha - k_{\alpha,\beta'}$, the $n_\alpha - k_{\alpha,\beta'}$ affine charts

$$(\mathbb{C}^{n_{\alpha'}})_{\alpha'}, \quad \text{where } \alpha' := (\alpha, i), \quad i = 1, \dots, n_\alpha - k_{\alpha,\beta'}, \quad n_{\alpha'} := 2n_\alpha - k_{\alpha,\beta'}$$

We also create a new collection of open subsets $U_{\alpha',\beta'} \subset (\mathbb{C}^{n_{\alpha'}})_{\alpha'}$ by taking the inverse images of $U_{\alpha,\beta'} \subset (\mathbb{C}^{n_\alpha})_\alpha$ under the morphisms $(\mathbb{C}^{n_{\alpha'}})_{\alpha'} \rightarrow (\mathbb{C}^{n_\alpha})_\alpha$.

The birational maps. The natural projection $\pi_\alpha : (\mathbb{C}^{n_{\alpha'}})_{\alpha'} \rightarrow (\mathbb{C}^{n_\alpha})_\alpha$ on the first n_α components defines the birational morphism $\pi_{\alpha',\beta'} = \pi_\alpha|_{X_{\alpha',\beta'}} : X_{\alpha',\beta'} \rightarrow X_{\alpha,\beta'}$, for any α, β' , such that $X_{\alpha,\beta'} \neq \emptyset$. This defines birational morphisms

$$i_{\alpha'_1\beta'_1, \alpha'_2\beta'_2} : X_{\alpha'_1\beta'_1} \xrightarrow{\pi_{\alpha'_1\beta'_1}^{-1}} X_{\alpha_1\beta'_1} \xrightarrow{i_{\alpha_1\beta'_1, \alpha'_2\beta'_2}} X_{\alpha_2\beta'_2} \xleftarrow{\pi_{\alpha'_2\beta'_2}^{-1}} X_{\alpha'_2\beta'_2}.$$

Consider a blow-up $X_{\mathcal{T}'}$ of $C \subset X_{\mathcal{T}}$. There exist open embeddings $j_{\alpha',\beta'} : X_{\alpha',\beta'} \cap U_{\alpha',\beta'} \hookrightarrow X_{\mathcal{T}'}$ induced by $j_{\alpha,\beta} : X_{\alpha,\beta} \cap U_{\alpha,\beta} \hookrightarrow X_{\mathcal{T}}$, defining an open cover of $X_{\mathcal{T}'}$, and satisfying

$$j_{\alpha'_2\beta'_2}^{-1} \circ j_{\alpha'_1\beta'_1} = i_{\alpha'_1\beta'_1, \alpha'_2\beta'_2}.$$

Equations of blow-up. Without loss of generality the blow-up in each of the $n_\alpha - k_{\alpha\beta'}$ affine charts $(\mathbb{C}^{n_{\alpha'}})_{\alpha'}$, where $\alpha' := (\alpha, i)$, $i = 1, \dots, n_\alpha - k_{\alpha\beta'}$, can be described as follows: (For simplicity, we drop the α, β indices below.)

Assume that the function u_{i_0} , $i_0 \leq n - k$, defines the chart of the blow-up. The blow-up of \mathbb{C}^n is a closed subset $bl(\mathbb{C}^n)$ of \mathbb{C}^{2n-k} described by the following equations

$$\begin{aligned} u_j - u_{i_0}x_{j+n} &= 0 \quad \text{for } 0 < j \leq n - k, j \neq i_0, \\ u_{i_0} - x_{i_0+n} &= 0. \end{aligned}$$

The exceptional divisors. The exceptional divisor for this blow-up is given by $u_{i_0} = 0$ on $bl(\mathbb{C}^n)$. Since $u_{i_0} = x_{i_0+n}$ we may represent it by the coordinate x_{i_0+n} on \mathbb{C}^{2n-k} .

The previous exceptional divisors keep their form $x_j = 0$ if they do not describe C , or they convert to $x_{j+n}(= u_j/u_{i_0})$ if they were described by the function $u_j \equiv x_j$.

The strict transform of X . Recall that X is described by $u_1 = \dots = u_{n-m} = 0$ on $U \subset \mathbb{C}^n$. The blow-up of $X = X_{\alpha,\beta}$ is a closed subset $X' \subset \mathbb{C}^{2n-k}$ which is described by a new set of equations:

1. $u_j - u_{i_0}x_{j+n} = 0$ for $0 < j \leq n - k, j \neq i_0$;
2. $u_{i_0} - x_{i_0+n} = 0$;
3. $x_{j+n} = 0$ for $0 < j \leq n - m, j \neq i_0$.

In some situations we consider additionally the induced equation

$$(4) \quad 1 = 0 \quad \text{if } 0 < i_0 \leq n - m.$$

(Note that the equations of the first two types describe the blow-up $bl(\mathbb{C}^n)$ of \mathbb{C}^n . The third and the fourth types of the equations $x_{j+n} = u_j/u_{i_0} = 0, j \neq i_0$ (or $1 = u_{i_0}/u_{i_0} = 0$) describe the strict transform of X inside $bl(\mathbb{C}^n)$. In the latter case if $0 < j = i_0 \leq n - m$ the strict transform is an empty set in the relevant chart. Still we shall keep the uniform description of the objects and their transformations, and do not eliminate any equations in the description of the empty set.)

4.2.1. The generators of $\mathcal{I}_{\alpha,\beta}$ after blow-up. We will not compute the controlled transforms of the generators of $\mathcal{I} (= \mathcal{I}_{\alpha,\beta})$ (over U) directly. Instead we modify them first. The generators g_i of \mathcal{I} satisfy, by Lemma 2.2.1, the condition

$$g_i \cdot f^{r_i} \in \mathcal{I}_C^\mu + \mathcal{I}_X,$$

where $V(f) = \mathbb{C}^n \setminus U \subset \mathbb{C}^n$ (we have dropped the indices α, β here).

For any generator g_i write

$$g_i \cdot f^{r_i} = \sum_{a_{n-m+1} + \dots + a_{n-k} = \mu} h_{(a_{n-m+1}, \dots, a_{n-k}), i} u_{n-m+1}^{a_{n-m+1}} \cdot \dots \cdot u_{n-k}^{a_{n-k}} + \sum_{j=1, \dots, m-n} h_{ij} u_j.$$

Set $\bar{a} := (a_{n-m+1}, \dots, a_{n-k}), \bar{u}^{\bar{a}} := u_{n-m+1}^{a_{n-m+1}} \cdot \dots \cdot u_{n-k}^{a_{n-k}}$. Then we can rewrite the above as

$$g_i \cdot f^{r_i} = \sum_{|\bar{a}| = \mu} h_{\bar{a}} \bar{u}^{\bar{a}} + \sum_{j=1, \dots, m-n} h_{ij} u_j.$$

To bound r_i and $\deg(h_{\bar{a},i}), \deg(h_{ij})$ we first consider a similar equality,

$$g_i \cdot f^{R_i} = \sum_{|\bar{a}| = \mu} H_{\bar{a}} \bar{u}^{\bar{a}} + \sum_{j=1, \dots, m-n} H_{ij} u_j,$$

for certain $R_i, H_{\bar{a},i}, H_{ij}$. We introduce a new variable z and we get an equality

$$g_i = z^{R_i} \cdot \left(\sum_{|\bar{a}|=\mu} H_{\bar{a}} \bar{u}^{\bar{a}} + \sum_{j:=1,\dots,m-n} H_{ij} u_j \right) + g_i \cdot \left(\sum_{0 \leq j \leq r-1} (z \cdot f)^j \right) \cdot (1 - z \cdot f);$$

in other words, g_i belongs to the ideal generated by $\{\bar{u}^{\bar{a}}\}, \{u_j\}, 1 - z \cdot f$. Therefore one can represent g_i as

$$g_i = \sum_{|\bar{a}|=\mu} \tilde{H}_{\bar{a},i} \cdot \bar{u}^{\bar{a}} + \sum_{j:=1,\dots,m-n} \tilde{H}_{ij} \cdot u_j + \tilde{H} \cdot (1 - z \cdot f),$$

for suitable polynomials $\tilde{H}_{\bar{a},i}, \tilde{H}_{ij}, \tilde{H}$ with degrees less than $(d \cdot \mu)^{2^{O(n)}}$ due to [45], [25], [42]. Hence substituting in the latter equality $z = 1/f$ and cleaning the denominator we obtain the bound $(d \cdot \mu)^{2^{O(n)}}$ on $r_i, \text{deg}(h_{\bar{a},i}), \text{deg}(h_{ij})$.

The generators after blow-up and their degree. Using the above we can describe the controlled transform of \mathcal{I} in terms of the controlled transforms of its modified generators. Define the modified generators of \mathcal{I} to be

$$\bar{g}_i = \sum_{|\bar{a}|=\mu} h_{\bar{a}} \bar{u}^{\bar{a}}.$$

Then their controlled transforms are given by

$$\sigma^c(\bar{g}_i) = u_{n-k}^{-\mu} \sigma^*(h_{\bar{a},i} u^{\bar{a}}).$$

Denote by $G(d, n, \mu)$ the bound on the degree of the resulting marked ideal \mathcal{T}' after a blow up applied to a marked ideal \mathcal{T} , provided that $d(\mathcal{T}) \leq d, n(\mathcal{T}) \leq n$. Thus, by the above:

LEMMA 4.2.1. $G(n, d, \mu) < (d \cdot \mu)^{2^{O(n)}}$.

4.3. Elementary operations and elementary auxillary functions. To estimate the complexity of the desingularization algorithm we introduce a few auxiliary functions related to the ingredients of \mathcal{T} . It is convenient to associate to \mathcal{T} with data (m, d, n, l, q, μ) , the vector

$$\gamma := (r, m, d, n, l, q, \mu) \in \mathbf{Z}_{\geq 0}^7,$$

where r is the subscript of the element of the resolution $(\mathcal{T}_r)_{r=0,1,\dots}$,

4.3.1. The effect of a single blow-up. Summarizing the above we get the following:

LEMMA 4.3.1. *Consider the object $\mathcal{T} := (\{X_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}, E_{\alpha,\beta}, U_{\alpha,\beta}\}_{\alpha,\beta}, n, \mu)$ with data (m, μ, d, n, l, q) . Let $\mathcal{T}' := (\{X'_{\alpha',\beta'}, \mathcal{I}'_{\alpha',\beta'}, E'_{\alpha',\beta'}, U'_{\alpha',\beta'}\}_{\alpha',\beta'}, n_{\alpha'}, \mu)$ denote the object with data (m, μ, d', n', l', q') obtained from \mathcal{T} by a single blow-up at the center C represented by the collection of closed sets $\{C_{\alpha\beta''} \subset (\mathbb{C}^{n_\alpha})_\alpha\}$ describing the center in open subsets $U_{\alpha\beta''} \subset U_{\alpha\beta}$. Assume that the maximal degree of the polynomials describing the center is less than d . Assume that q gives also a bound for the number of open neighborhoods $U_{\alpha\beta''}$. Then*

1. $d' \leq G(n, d, \mu) < (d \cdot \mu)^{2^{O(n)}}$;
2. $n' \leq 2n$;

3. $l' \leq l + n$;
4. $q' \leq n \cdot q$;

The effect of the single blow-up can be measured by the function

$$Bl(r, m, d, n, l, q, \mu) := (r + 1, m, G(n, d, \mu), 2n, l + n, n \cdot q, \mu).$$

The multiple effect of the t blow-ups can be measured by the recursive function

$$\overline{Bl}(r, m, d, n, l, q, \mu, t) = Bl \circ \overline{Bl}(r, m, d, n, l, q, \mu, t - 1),$$

or, more briefly,

$$\overline{Bl}(\gamma, t) = Bl(\overline{Bl}(\gamma, t - 1)).$$

Note that

$$\overline{Bl}(r, m, d, n, l, q, \mu, t) = (r + t, m, \overline{G}(n, d, \mu, t), 2^t n, l + 2^{t-1} n, (2^{t+1} - 1) \cdot n^t q, \mu)$$

for the relevant function $\overline{G}(n, d, \mu, t)$.

5. Bounds on multiplicities and degrees of coefficient ideals.

5.0.2. The maximal multiplicity of \mathcal{I} on the subvariety $X \subset \mathbb{C}^n$. Let d be the maximal degree of (X, \mathcal{I}) on \mathbb{C}^n . Denote by $M(d, n)$ a bound on the multiplicity of the ideal \mathcal{I} on X . To estimate $M(d, n)$, we can assume (after a linear transformation of the coordinates) that the order of the polynomial $u_{i_j} - x_j$ is at least 2 for all $1 \leq j \leq n - m$. For any polynomial $g \in \mathcal{I}_{\alpha, \beta}$ one can find polynomials $h_j \in \mathbb{C}[x_1, \dots, x_n]$, $0 \leq j \leq n - m$ and $h \in \mathbb{C}[x_{n-m+1}, \dots, x_n]$ such that

$$(5) \quad h_0 \cdot g + \sum_{1 \leq j \leq n-m} h_j \cdot u_{i_j} = h(x_{n-m+1}, \dots, x_n).$$

Moreover, we can rewrite the latter equality over the field $\mathbb{C}(x_{n-m+1}, \dots, x_n)$ in the form

$$(6) \quad \tilde{h}_0 \cdot g + \sum_{1 \leq j \leq n-m} \tilde{h}_j \cdot u_{i_j} = 1,$$

where $\tilde{h}_j = \frac{h_j}{h} \in \mathbb{C}(x_{n-m+1}, \dots, x_n)[x_1, \dots, x_{n-m}]$, with the common denominator in $\mathbb{C}[x_{n-m+1}, \dots, x_n]$ for $0 \leq j \leq n - m$.

We apply to (6) the Effective Nullstellensatz (see e.g. [16], [24], [37], [35]). This gives us the bound $d^{O(n)}$ on the degrees of $\tilde{h}_i = h_i/h$ with respect to variables x_1, \dots, x_{n-m} (for certain solutions). To find h_i/h one can solve the latter equality considering it as a linear system over the field $\mathbb{C}(x_{n-m+1}, \dots, x_n)$.

The algorithm can find

$$\tilde{h}_j = \sum a_{I,j} x^I$$

with indeterminates $a_{I,j} \in \mathbb{C}(x_{n-m+1}, \dots, x_n)$, and monomials $x^I = x_1^{i_1} \cdot \dots \cdot x_{n-m}^{i_{n-m}}$ with degrees $i_1 + \dots + i_{n-m} \leq d^{O(n)}$ substituting \tilde{h}_j in (6) and solving linear system over the field $\mathbb{C}(x_{n-m+1}, \dots, x_n)$. Clearing the common denominator in (6) gives (5) with

$$\deg(h), \deg(h_j) \leq d^{O(n^2)}, \quad 1 \leq j \leq n - m$$

Now we have to estimate the maximal multiplicity $\text{ord}_x(g|_X)$ for $g \in \mathbb{C}[x_1, \dots, x_n]$, such that $\deg(g) \leq d$ and $x \in X$. We use (5). We get immediately by the above:

LEMMA 5.0.2. *The maximal multiplicity $\text{ord}_x(g|_X)$ on X for any function $g \in \mathbb{C}[x_1, \dots, x_n]$, such that $\deg(g) \leq d$ and $x \in X$, is bounded by the function $M(d, n) = d^{O(n^2)}$ constructed as above:*

$$\text{ord}_x(g|_X) \leq M(d, n).$$

Proof. $\text{ord}_x(g|_X) \leq \text{ord}_x(h_0 \cdot g)|_X = \text{ord}_x h(x_{n-m+1}, \dots, x_n)|_X = \text{ord}_x h(x_{n-m+1}, \dots, x_n) \leq \deg(h) \leq M(d, n) = d^{O(n^2)}$. \square

5.0.3. Derivations on the subvariety $X \subset \mathbb{C}^n$. In order to follow the construction of the algorithm we use the language of derivations Der_X on X . Since our X is embedded into \mathbb{C}^n it is natural to represent all objects on X as the restriction of the relevant objects on \mathbb{C}^n to $X \subset \mathbb{C}^n$. Unfortunately the sheaf of derivations on \mathbb{C}^n does not restrict well to X :

$$\text{Der}_{\mathbb{C}^n|_X} \neq \text{Der}_X.$$

Instead we consider

$$\text{Der}_{\mathbb{C}^n, X} := \{D \in \text{Der}_{\mathbb{C}^n} \mid D(\mathcal{I}_X) \subset (\mathcal{I}_X)\}.$$

LEMMA 5.0.3. *Let $u_1 = 0 \dots, u_{n-m} = 0$ describe $X \subset \mathbb{C}^n$. Then the ring $\text{Der}_{\mathbb{C}^n, X}$ is generated by*

$$\{u_i \cdot d_{u_j}\}_{1 \leq i \leq n-m, 1 \leq j \leq n-m} \cup \{d_{u_j}\}_{n-m < j \leq n}.$$

In particular the restriction $\text{Der}_{\mathbb{C}^n, X}|_X = \text{Der}_X$ is generated by

$$\{d_{u_j}\}_{n-m < j \leq n}.$$

Proof. This follows from the definition. \square

Note that since we have

$$(d_{u_j})_j = ((\partial u_j / \partial x_i)_{i,j})^{-1} \cdot (d_{x_i})_i = \frac{1}{\det((\partial u_j / \partial x_i)_{i,j})} \text{adj}((\partial u_j / \partial x_i)_{i,j}) \cdot (d_{x_i})_i$$

we immediately get the following:

LEMMA 5.0.4. *Let $\mathcal{U} := \{x \in \mathbb{C}^n \mid \det((\partial u_j / \partial x_i)_{i,j}) \neq 0\}$. The sheaf $\text{Der}_{\mathbb{C}^n, X}$ is generated over $U \subset \mathbb{C}^n$ by*

$$(7) \quad \{u_i \cdot d'_{u_j}\}_{1 \leq i \leq n-m, 1 \leq j \leq n-m} \cup \{d'_{u_j}\}_{n-m < j \leq n},$$

where

$$(d'_{u_j})_j = \text{adj}((\partial u_j / \partial x_i)_{i,j}) \cdot (d_{x_i})_i.$$

The derivations (7) generate a subsheaf $\overline{\text{Der}_{\mathbb{C}^n, X}}$ of $\text{Der}_{\mathbb{C}^n, X}$ over \mathbb{C}^n . Both sheaves coincide over U . Thus we will replace $\text{Der}_{\mathbb{C}^n, X}$ with $\overline{\text{Der}_{\mathbb{C}^n, X}}$ for our computations over U .

LEMMA 5.0.5. *Let \mathcal{I} be any ideal on \mathbb{C}^n . Assume the maximal degree of some generating set of \mathcal{I} is $\leq d_1$, and the maximal degree of u_i is less than d_2 . Then the maximal degree of generators of $\text{Der}_{\mathbb{C}^n, X}(\mathcal{I})$ is bounded by $d_1 + nd_2$.*

5.0.4. Construction of the coefficient homogenized companion ideal.

Recall that in step 2 for the marked ideal (\mathcal{I}, μ) , we find the maximal multiplicity $\bar{\mu} \leq M(d, n)$ of the nonmonomial part $\mathcal{N}(\mathcal{I})$, and construct the companion ideal $\mathcal{O}(\mathcal{I})$, for which we immediately take homogenized coefficient ideal $\mathcal{J} := \mathcal{C}(\mathcal{H}(\mathcal{O}(\mathcal{I})))$. In our situation of the set $X \subset \mathbb{C}^n$ defined by set of parameters $\{u_i\}_{1 \leq i \leq n-m}$ on the open set $U \subset \mathbb{C}^n$ we will use $\overline{\text{Der}}_{\mathbb{C}^n, X}$ instead of Der_X for the definition above. Immediately from the definition, we get a formula for a bound $A(d, n, \mu)$ on the degrees of generators of the marked ideal $\mathcal{C}(\mathcal{H}(\mathcal{I}_q))$. Note first that we have the the following bounds on the multiplicities:

LEMMA 5.0.6. $\mu(\mathcal{N}(\mathcal{I})) = \bar{\mu}$, $\mu(\mathcal{O}(\mathcal{I})) \leq \mu \cdot \bar{\mu}$ and $\mu(\mathcal{J}) \leq (\mu \cdot \bar{\mu})! \leq (\mu \cdot M(d, n))!$.

As a Corollary, we obtain:

LEMMA 5.0.7. *The maximal degree of generators of $\mathcal{J} := \mathcal{C}(\mathcal{H}(\mathcal{O}(\mathcal{I})))$ is bounded by*

$$A(d, n, \mu) := (\mu \cdot \bar{\mu})!nd \leq (\mu \cdot M(d, n))!nd \leq (d^{\mathcal{O}(n^2)})!$$

Proof. This follows from Lemma 5.0.5. \square

5.0.5. Restriction to hypersurfaces of maximal contact, and to exceptional divisors. In step 1, we restrict \mathcal{I} to intersections of the exceptional divisors and maximal contact.

We need to estimate a bound $B(d, n, \mu)$ for the degree of the maximal contact $u \in \overline{\text{Der}}_{\mathbb{C}^n, X}^{\bar{\mu}-1}(\mathcal{N}(\mathcal{I}))$. The following is an immediate consequence of Lemma 5.0.5:

LEMMA 5.0.8. *The maximal degree of any maximal contact $u \in \overline{\text{Der}}_{\mathbb{C}^n, X}^{\bar{\mu}-1}(\mathcal{N}(\mathcal{I}))$ is bounded by*

$$B(d, n, \mu) := \bar{\mu}nd \leq M(d, n) \cdot nd \leq d^{\mathcal{O}(n^2)}.$$

5.0.6. A bound for the number of generators of J . First, we state the basic properties concerning the number of generators of the ideal in the following lemma:

LEMMA 5.0.9.

1. *The number of generators of $\overline{\text{Der}}_{\mathbb{C}^n, X}(\mathcal{I})$ is given by $(n+1)l(\mathcal{I})$.*
2. *The number of generators of $\overline{\text{Der}}_{\mathbb{C}^n, X}^i(\mathcal{I})$ is given $(n+1)^i l(\mathcal{I})$.*
3. *The number of generators of \mathcal{I}^i is bounded by $l(\mathcal{I})^i$.*
4. *The number of generators of $\mathcal{I}' = \mathcal{O}(\mathcal{I})$ can be bounded by $L_{\mathcal{O}}(l(\mathcal{I}), \mu) := l(\mathcal{I})^\mu + 1$.*
5. *The number of generators of $\mathcal{I}' = \mathcal{H}(\mathcal{I}, \mu)$ can be bounded by $L_{\mathcal{H}}(l(\mathcal{I}), \mu) := \mu(n+1)^{\mu^2} l(\mathcal{I})^\mu$.*
6. *The number of generators of $\mathcal{I}' = \mathcal{C}(\mathcal{I}, \mu)$ can be bounded by $L_{\mathcal{C}}(l(\mathcal{I}), \mu) := \mu(n+1)^{\mu^1} l(\mathcal{I})^{\mu^1}$.*

Proof. Immediate from the definitions. \square

By the preceding lemma, we get

$$\begin{aligned} l(\mathcal{H}(\mathcal{O}(\mathcal{I}))) &\leq L_{\mathcal{H}}(L_{\mathcal{O}}(l(\mathcal{I}), \mu), \mu \cdot M(d, n)), \\ l(\mathcal{J}) = l(\mathcal{C}(\mathcal{H}(\mathcal{O}(\mathcal{I})))) &\leq L_{\mathcal{C}}(L_{\mathcal{H}}(L_{\mathcal{O}}(l(\mathcal{I}), \mu), \mu \cdot M(d, n)), \mu \cdot M(d, n)). \end{aligned}$$

Thus we get:

COROLLARY 5.0.10. $l(\mathcal{J}) \leq F(d, n, \mu)$, where

$$F(d, n, \mu) := L_C(L_{\mathcal{H}}(L_{\mathcal{O}}(l(\mathcal{I}), \mu), \mu \cdot M(d, n)), \mu \cdot M(d, n)).$$

REMARK. The algorithm of [12] does not involve the homogenization step and therefore gives better estimates for the elementary functions introduced. In particular:

1. The degree of generators of \mathcal{J} is bounded by $B(d, n, \mu)$ (which improves the bound $A(d, n, \mu)$, cf. Lemma 5.0.7).
2. The number of generators $l(\mathcal{J})$ can be bounded by $L_C(L_{\mathcal{O}}(l(\mathcal{I}), \mu), \mu \cdot M(d, n))$ (which improves the bound $F(d, n, \mu)$; cf. Corollary 5.0.10),

However, the above improvements do not affect the overall Grzegorzczuk complexity class \mathcal{E}^{m+3} . (See Theorem 6.4.2.)

Summarizing:

LEMMA 5.0.11. *The effect of passing from \mathcal{I} to $\mathcal{J} = (\mathcal{C}(\mathcal{H}(\mathcal{O}(\mathcal{I}))))$ as in Step 2a/Step1 can be described by the function*

$$\Delta_{2a}(r, m, d, n, l, q, \mu) := (r, m, A(d, n, \mu), n, F(d, n, \mu, l), q, (\mu \cdot M(d, n))!).$$

5.0.7. A bound for the number of maximal contacts and the relevant neighborhoods. We will construct maximal contacts along with the open neighborhoods for which they are defined. Each maximal contact $u \in \overline{\text{Der}_{\mathbb{C}^n, X}^{\bar{\mu}-1}}(\mathcal{N}(\mathcal{I}))$ that we consider is of the form $u = D^{\bar{\mu}-1}(g_i)$, where $D^{\bar{\mu}-1} = D_1^{a_1} \dots D_n^{a_n}$ is a certain composition of $\bar{\mu} - 1$ differential operators (7) (i.e., the D_i are of the form in (7), and $a_1 + \dots + a_n = \bar{\mu} - 1$). Consider all differential operators $\{D_r^{\bar{\mu}}\}_{r \in R}$, which are certain compositions of $\bar{\mu}$ differential operators (7), and take all the corresponding functions $f_{r,i} := D_r^{\bar{\mu}}(g_i)$. On the open set $U_{r,i} = U \setminus V(f_{r,i})$, consider the maximal contact $u_{r,i} = D^{\bar{\mu}-1}(g_i)$, where $D^{\bar{\mu}-1} = D_1^{a_1} \dots D_n^{a_n}$ is a certain composition of $\bar{\mu} - 1$ differential operators (7) obtained from $D_r^{\bar{\mu}} = D_1^{b_1} \dots D_n^{b_n}$ by replacing one of the positive b_i with $b_i - 1$ (i.e. $a_i := b_i - 1$ for some $b_i > 0$ and $a_j = b_j$ for $j \neq i$.)

LEMMA 5.0.12. *The number of the maximal contacts $u_{r,i} \in \overline{\text{Der}_{\mathbb{C}^n, X}^{\bar{\mu}-1}}(\mathcal{N}(\mathcal{I}))$ and at the same time the number of neighborhoods $U_{i,r} \subset U$ can be bounded by*

$$C(d, n, \mu) \cdot l(\mathcal{I}),$$

where $l(\mathcal{I})$ is the number of generators of \mathcal{I} , and

$$C(d, n, \mu) := \binom{M(d, n) + n}{n}.$$

Proof. The number of the maximal contacts is bounded by $\binom{\bar{\mu}+n}{n} \cdot l(\mathcal{I}) \leq C(d, n, \mu) \cdot l(\mathcal{I})$, where $C(d, n, \mu) := \binom{M(d, n)+n}{n}$. \square

Summarizing:

LEMMA 5.0.13. *The effect of passing from \mathcal{I} to $\mathcal{J} = (\mathcal{C}(\mathcal{H}(\mathcal{O}(\mathcal{I}))))$ and then to \mathcal{J}_{H_s} in Step 1a or to $\mathcal{J}_{|V(u)}$ as in the Step 1b can be described by the function*

$$\Delta_1(r, m, d, n, l, q, \mu) := (r, m-1, A(d, n, \mu), n, F(d, n, \mu, l), q \cdot l \cdot C(d, n, \mu), (\mu \cdot M(d, n))!).$$

Note that the restriction to the maximal contact does not affect the degree since the function $B(d, n, \mu)$ measuring the degree of maximal contact is smaller than $A(d, n, \mu)$.

REMARK. The particular form of the bounds obtained does not strongly influence Theorem 6.4.2; we need only that the functions belong to the class \mathcal{E}^3 . (See the beginning of the next section.)

6. Complexity bound of the resolution algorithm in terms of Grzegorzczuk's classes.

6.1. Language of Grzegorzczuk's classes. The complexity estimate of the desingularization algorithm which we provide in this section is given in terms of Grzegorzczuk's classes \mathcal{E}^l , $l \geq 0$ of primitive-recursive functions [28], [52]. To make the paper self-contained, we provide a definition of \mathcal{E}^l by induction on l (informally speaking, \mathcal{E}^l consists of integer functions $\mathbb{Z}^s \rightarrow \mathbb{Z}^t$ whose construction requires l nested primitive recursions).

For the base definition, \mathcal{E}^0 contains constant functions $x_k \mapsto c$, functions $x_k \mapsto x_k + c$ and projections $(x_1, \dots, x_n) \mapsto x_k$ for any variables x_1, \dots, x_n .

The class \mathcal{E}^1 contains linear functions $x_k \mapsto c \cdot x_k$ and $(x_{k_1}, x_{k_2}) \mapsto x_{k_1} + x_{k_2}$.

The class \mathcal{E}^2 contains all polynomials with integer coefficients

Let $l \geq 2$. For the inductive step of the definition, assume that functions $G(x_1, \dots, x_n), H(x_1, \dots, x_n, y, z) \in \mathcal{E}^l$. Then the function $F(x_1, \dots, x_n, y)$ defined by the primitive recursion,

$$(8) \quad F(x_1, \dots, x_n, 0) = G(x_1, \dots, x_n),$$

$$(9) \quad F(x_1, \dots, x_n, y + 1) = H(x_1, \dots, x_n, y, F(x_1, \dots, x_n, y)),$$

belongs to \mathcal{E}^{l+1} .

To complete the definition of \mathcal{E}^l , $l \geq 0$, take the closure with respect to composition and the following *limited primitive recursion*:

Let $G(x_1, \dots, x_n), H(x_1, \dots, x_n, y, z), Q(x_1, \dots, x_n, y) \in \mathcal{E}^l$. Then the function $F(x_1, \dots, x_n, y)$ defined by (8),(9) also belongs to \mathcal{E}^l , *provided that*

$$F(x_1, \dots, x_n, y) \leq Q(x_1, \dots, x_n, y).$$

Clearly, $\mathcal{E}^{l+1} \supset \mathcal{E}^l$.

Observe that \mathcal{E}^3 contains all towers of exponential functions and \mathcal{E}^4 contains all tetration functions [28], [52].

The union $\cup_{l < \infty} \mathcal{E}^l$ coincides with the set of all primitive-recursive functions.

6.2. Resolution algorithm as a graph. It is instructive to represent the resolution algorithm in the form of a tree T as in the following Figure.

Each node a of T corresponds to an intermediate object $T_a = (\{X_{\alpha,\beta}, \mathcal{I}_{\alpha,\beta}, E_{\alpha,\beta}, U_{\alpha,\beta}\}_{\alpha,\beta}, n, \mu)$. Each node a is labeled either by 1 or 2 depending on whether it corresponds to step 1 or 2 in the description of the algorithm (see the previous sections). An edge from a node labeled by 1 leads to its child node labeled by 2 and the edge is labeled in its turn either by 2a or by 2b depending on the step to which it corresponds. Similarly, an edge from a node labeled by 2 leads to its child node labeled by 1 and is labeled in turn either by 1a or by 1b. In the Figure a child node is always located to the right from a node.

The algorithm yields T by recursion starting with its root. Assume that a and T_a are already constructed. The next task of the algorithm is to resolve the object

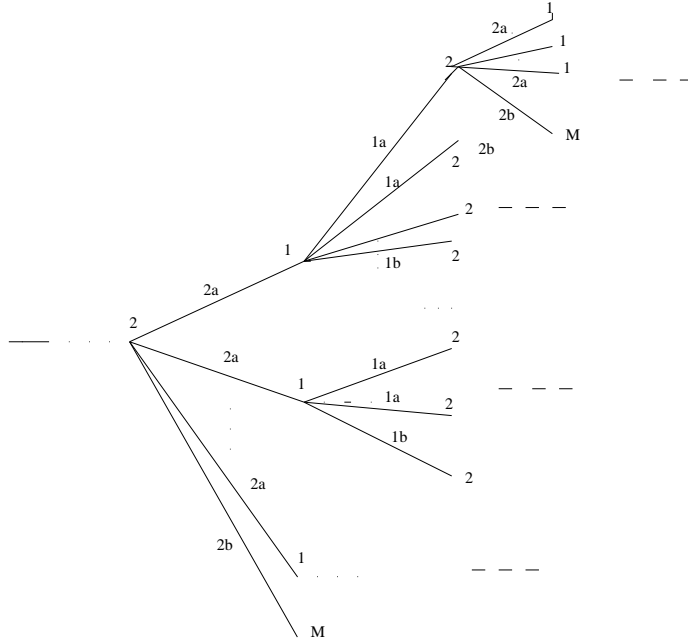


FIG. 1.

T_a . To this end the algorithm first constructs the child nodes $a_1, \dots, a_{\bar{t}}$ of a according to the algorithm. The order of producing $a_1, \dots, a_{\bar{t}}$ goes from up to down in the Figure. The algorithm resolves the objects $T_{a_1}, \dots, T_{a_{t-1}}$ by recursion on $0 \leq t \leq \bar{t}$ and in the process modifies $T_a := T_a(0)$, obtaining the current object $T_a(t-1)$. Then the algorithm yields a_t, T_{a_t} , resolves T_{a_t} and collects all the blow ups produced while resolving T_{a_t} and applies them (with the same centers) to the current object $T_a(t-1)$; the resulting object we denote by $T_a(t)$. This allows the algorithm to yield a child node a_{t+1} and $T_{a_{t+1}}$ following the description from the previous sections.

For the leaves of T there are two possibilities: either a leaf is labeled by 2 or a certain node a labeled by 2 could have a single edge (the lowest in the Figure among the edges originating at a) labeled by 2b which leads to a child node of a being a leaf corresponding to the monomial case (labeled by M). Note also that if a is labeled by 1 then the top few edges originating at a are labeled by 1a and the remaining bottom ones are labeled by 1b (in the order from up to down in the Figure).

Observe that the dimension of the varieties $X_{\alpha,\beta}$ corresponding to a drops while passing to any of its child nodes when a is labeled by 1, and the dimension does not increase when a is labeled by 2. Therefore, the depth of T does not exceed $2m$.

6.3. Main recursive functions. Now we proceed to the bounds of some recursive functions related to the ingredients of \mathcal{T} . Set

$$\gamma := (r, m, d, n, l, q, \mu) \in \mathbf{Z}_{\geq 0}^7.$$

Let \mathcal{T}_* be the canonical resolution of \mathcal{T} . For simplicity of notation, we introduce the following function defined on $\mathbf{Z}_{\geq 0}^7$:

$$\Gamma^{(m)}(\gamma) := (r + R^{(m)}(\gamma), m, D^{(m)}(\gamma), N^{(m)}(\gamma), L^{(m)}(\gamma), Q^{(m)}(\gamma), \mu),$$

where the subscript r can be interpreted as the subscript in the resolution of \mathcal{T} , and

1. $R^{(m)}(\gamma)$ is the number of blow-ups needed to resolve the initial marked ideal with data bounded by (m, d, n, l, q, μ) ;
2. $D^{(m)}(\gamma)$ is a function bounding the maximum of the degrees of all the polynomials which represent \mathcal{T}_* and all objects constructed along the way (in particular, the centers);
3. $N^{(m)}(\gamma)$ is a bound for the dimensions of the ambient affine spaces constructed along the way;
4. $L^{(m)}(\gamma)$ is a bound for the number of polynomials appearing in the description of a single neighborhood $U_{\alpha, \beta}$ on resolving T ;
5. $Q^{(m)}(\gamma)$ is the number of neighborhoods in all the auxiliary objects (in particular, the centers) appearing on resolving T .

REMARK. The functions $R^{(m)}(\gamma)$, $D^{(m)}(\gamma)$, $N^{(m)}(\gamma)$ do not depend on l, q .

6.3.1. Algorithm revisited. Let (\mathcal{I}, μ) be a marked ideal on an m -dimensional smooth variety X . Consider the corresponding object

$$\mathcal{T}^{(m)} = (\{X_{\alpha, \beta}, \mathcal{I}_{\alpha, \beta}, E_{\alpha, \beta}, U_{\alpha, \beta}, \mathbb{C}_\alpha^n \mid \alpha \in A, \beta \in B\}, \mu)$$

with the initial data

$$\gamma := (0, m, d, n, l, q, \mu).$$

Our next goal is two-fold. We will give recursive formulas for $\Gamma^{(m)}, R^{(m)}, D^{(m)}, N^{(m)}, L^{(m)}, Q^{(m)}$ and prove by induction on m these functions belong to Grzegorzczuk's class \mathcal{E}^{m+3} . In the base of the induction (i.e., for $m = 0$), the functions

$$R^{(0)} = 1, D^{(0)} = O(dn), N^{(0)} \leq 2n, L^{(0)} \leq l \cdot (dn)^{O(n)}, Q^{(0)} \leq nq$$

belong to the class \mathcal{E}^3 , as does $\Gamma^{(0)}$.

Now we proceed to the inductive step. Assume that $\Gamma^{(m-1)}, R^{(m-1)}, D^{(m-1)}, N^{(m-1)}, L^{(m-1)}, Q^{(m-1)}$ belong to Grzegorzczuk's class \mathcal{E}^{m+2} , where $m \geq 1$.

If $\mathcal{I} = 0$, then the resolution is done by the single blow-up at the center $C = X$ and the object $\mathcal{T}^{(m)}$ is transformed into an object $\mathcal{T}_1^{(m)}$ with $X_1 = \emptyset$ and data bounded by $Bl(\gamma) \in \mathcal{E}^3$ (see Lemma 4.3.1).

If $\mathcal{I} \neq 0$, then the resolution algorithm can be represented by the following scheme.

Step 2. Resolve (\mathcal{I}, μ) on the m -dimensional smooth variety X . Consider the corresponding object $\mathcal{T}^{(m)}$ with initial data $\gamma := (0, m, d, n, l, q, \mu)$. Let $\bar{\mu}$ denote the maximal order of $\mathcal{N}(\mathcal{I})$ on X . We have the following estimate for $\bar{\mu}$:

$$\bar{\mu} \leq M(d, n) \in \mathcal{E}^3$$

(cf Lemma 5.0.2).

Step 2a. In this Step we are going to decrease the maximal order of the non-monomial part $\mathcal{N}(\mathcal{I})$ by resolving the companion ideal $O(\mathcal{I}, \mu)$. In fact we perform an additional modification of $\mathcal{O}(\mathcal{I})$ and construct the ideal $\mathcal{J} := \mathcal{C}(\mathcal{H}(\mathcal{O}(\mathcal{I})))$. This corresponds to the new object $\mathcal{T}_1^{(m)}$ with the initial data

$$\gamma^{(2a)} := \Delta_{2a}(\gamma) \in \mathcal{E}^3,$$

(see Lemma 5.0.11).

The object $\mathcal{T}_1^{(m)}$ will be then resolved and its resolution will cause the maximal order to decrease. The resolution $\mathcal{T}_1^{(m)}$ is done by performing Step 1.

Step 1. In this step we resolve \mathcal{J} , i.e., $\mathcal{T}_1^{(m)}$. The Step splits into two Steps (1a) and (1b).

Step 1a. Move apart all unions of the intersections $H_\alpha^s \subset \mathbb{C}_\alpha^n$ of s divisors in E , where s is the maximal number of divisors in E through points in $\text{supp } \mathcal{J} = \text{supp}(\mathcal{J}, \mu(\mathcal{J}))$. For any α , resolve all $\mathcal{J}_{|H_\alpha^s}$. We construct a new object

$$\mathcal{T}_2^{(s)} := (\{H_{\alpha,\beta}^s, \mathcal{J}_{\alpha,\beta}, E_{\alpha,\beta}, U_{\alpha,\beta}, \mathbb{C}_\alpha^n \mid \alpha \in A, \beta \in B\}, \mu(\mathcal{J})),$$

with initial data bounded by

$$\gamma^{(1a)} := \Delta_1(\gamma) \in \mathcal{E}^3$$

with $s \leq m - 1$ (see Lemma 5.0.13).

By the inductive assumption, the resolution of $\mathcal{T}_2^{(s)}$, i.e., the sequence $\mathcal{T}_{2*}^{(s)}$ of the induced intermediate objects determined by the blow-ups, requires at most $R^{(m-1)}(\gamma^{(1)})$ blow ups. The maximal degree of the polynomials of the centers and the objects $\mathcal{T}_{2*}^{(s)}$ describing the resolution does not exceed $D^{(m-1)}(\gamma^{(1)})$. The dimension n of the objects does not exceed $N^{(m-1)}(\gamma^{(1)})$.

Note that the resolution of $\mathcal{T}_{2*}^{(s)}$ determines a multiple blow-up $\mathcal{T}_{1*}^{(m)}$ of $\mathcal{T}_1^{(m)}$ consisting of $R^{(m-1)}(\gamma^{(1)})$ blow-ups. We have a direct correspondence between objects $\mathcal{T}_{1*}^{(m)}$, and $\mathcal{T}_{2*}^{(s)}$. The bound

$$\Gamma^{(m-1)}(\Delta_1(\gamma)) \in \mathcal{E}^{m+2},$$

for the data for the resolution $\mathcal{T}_{2*}^{(s)}$, given by the induction, remains valid for the data for $\mathcal{T}_{1*}^{(m)}$ as we use the same centers, the same ambient affine spaces, etc., for these multiple blow-ups. Only the strict transforms of the current X are different, and this does not affect the bounds for the data. We have additional equations to describe the current X in $\mathcal{T}_{2*}^{(s)}$, as compared to those in $\mathcal{T}_{1*}^{(m)}$.

Step 1a is performed at most $s \leq m$ times. Introduce the auxillary unknown $t = 0, 1, \dots, m$, and the function $\Gamma_{1a}^{(m)}(\gamma, t)$ which measures the possible effect after performing Step 1a t times:

$$\begin{aligned} \Gamma_{1a}^{(m)}(\gamma, 0) &:= \Delta_1(\gamma) \in \mathcal{E}^3, \\ \Gamma_{1a}^{(m)}(\gamma, t + 1) &:= \Gamma^{(m-1)}(\Gamma_{1a}^{(m)}(\gamma, t)). \end{aligned}$$

Since the Step 1a is performed at most m times, its final effect after completing Step 1a and passing to Step 1b is measured by the function

$$\Gamma_{1b}^{(m)}(\gamma) := \Gamma_{1a}^{(m)}(\gamma, m).$$

Note that for any fixed value of $t = t_0$ (in particular, for $t = m$) the functions $\Gamma_{1a}^{(m)}(\gamma, t_0)$, belong to the class \mathcal{E}^{m+2} due to the inductive hypothesis and Lemma 4.3.1, Corollary 5.0.10. Therefore, $\Gamma_{1b}^{(m)}$ belongs to the class \mathcal{E}^{m+2} . (We use here the property that Grzegorzcyk classes are closed under composition.) After performing Step 1a, we have moved apart all strict transforms of E and $\text{supp } \mathcal{J} = \text{supp}(\mathcal{J}, \mu(\mathcal{J}))$.

Step 1b. If the strict transforms of E do not intersect $\text{supp } \mathcal{J}$, we resolve \mathcal{J} , i.e., the object $\mathcal{T}_1^{(m)}$. This is achieved by resolving $\mathcal{J}|_{V(u)}$ (by induction), where $V(u)$ is a hypersurface of maximal contact. After completing Step 1a, the bound $\gamma^{(1a)}$ is transformed to

$$\gamma^{(1b)} := \Gamma_{1b}^{(m)}(\gamma) = (r^{(1b)}, m, d^{(1b)}, n^{(1b)}, l^{(1b)}, q^{(1b)}, (\mu \cdot (M(d, n))))!$$

(cf. Lemma 5.0.6).

Passing from \mathcal{J} to $\mathcal{J}|_{V(u)}$, we adjoin the equations of maximal contact as well as create new neighborhoods. This operation has been reflected in the construction of $\Delta_1(\gamma)$. By the construction of $\Gamma_{1b}^{(m)}(\gamma)$ and $\Delta_1(\gamma)$, the degree of the maximal contact does not exceed $d^{(1b)}$, while the number of neighborhoods does not exceed $q^{(1b)}$. In other words $\Gamma_{1b}^{(m)}(\gamma)$ bounds the initial data for $\mathcal{J}|_{V(u)}$.

The resolution process for $\mathcal{J}|_{V(u)}$ leads eventually to resolution of the object $\mathcal{T}_1^{(m)}$ corresponding to \mathcal{J} , with the data bounded by

$$\begin{aligned} \Gamma_1^{(m)}(\gamma) &:= \Gamma^{(m-1)}(\Gamma_{1b}^{(m)}(\gamma)) = \Gamma_{1a}^{(m)}(\gamma, m + 1) \\ &= (r^{(1)}, m, d^{(1)}, n^{(1)}, l^{(1)}, q^{(1)}, (\mu \cdot (M(d, n))))!, \end{aligned}$$

for the relevant $r^{(1)}, d^{(1)}, n^{(1)}, l^{(1)}, q^{(1)}$. Hence the function $\Gamma_1^{(m)}$ belongs to class \mathcal{E}^{m+2} by the inductive hypothesis and Lemma 4.3.1, Corollary 5.0.10. This completes Step 1.

The object $\mathcal{T}^{(m)}$ corresponding to \mathcal{I} with initial data γ is transformed to the new object with the data bounded by

$$\Gamma_{2a}^{(m)}(\gamma) := (r^{(1)}, m, d^{(1)}, n^{(1)}, l^{(1)}, q^{(1)}, \bar{\mu})$$

with smaller $\bar{\mu}$ — the maximal order of the new $\mathcal{N}(\mathcal{I})$. (Note that $\bar{\mu} < M(d, n)$).

This completes Step 2a. This Step 2a is then repeated at most $M(d, n)$ times until the maximal order drops to zero when we arrived at the monomial case. The final effect of Step 2a is measured then by the recursive function

$$\begin{aligned} \Gamma_{2a}^{(m)}(\gamma, 0) &= \gamma, \\ \Gamma_{2a}^{(m)}(\gamma, t + 1) &= \Gamma_{2a}^{(m)}(\Gamma_{2a}^{(m)}(\gamma, t)). \end{aligned}$$

Therefore, the function $\Gamma_{2a}^{(m)}$ belongs to class \mathcal{E}^{m+3} , by the definition of Grzegorzczuk classes (see (8),(9)), and by Lemmas 4.3.1, 5.0.2.

Putting $t = M(d, n)$ gives the final effect after completing all necessary Steps 2a and subsequently passing to Step 2b:

$$\Gamma_{2b}^{(m)}(\gamma) := \Gamma_{2a}^{(m)}(\gamma, M(d, n));$$

thereby, the function $\Gamma_{2b}^{(m)}$ belongs to class \mathcal{E}^{m+3} as well.

The procedure eventually reduces (\mathcal{I}, μ) to the monomial marked ideal $\mathcal{I} = \mathcal{M}(\mathcal{I})$.

Step 2b. Resolve the monomial marked ideal $\mathcal{I} = \mathcal{M}(\mathcal{I})$. The marked ideal corresponds to the object $\mathcal{T}^{(m)}$ with data

$$(r^{(2b)}, m, d^{(2b)}, n^{(2b)}, l^{(2b)}, q^{(2b)}, \mu) =: \Gamma_{2b}^{(m)}(\gamma).$$

The resolution of $\mathcal{I} = (x^\alpha)$ consists of blow-ups each of which decreases the multiplicity $|x^\alpha| \leq d^{(2b)}$. The resolution of \mathcal{I} requires at most $d^{(2b)}$ blow-ups. Thus the final solution data can be bounded by the function

$$\Gamma^{(m)}(\gamma) := \overline{Bl}(\Gamma_{2b}^{(m)}(\gamma), d^{(2b)}) \in \mathcal{E}^{m+3}.$$

We summarize the bounds achieved in the following Corollary (recall that the notation can be found in subsection 4.1.1).

COROLLARY 6.3.1. *When resolving a marked ideal (X, \mathcal{I}, E, μ) on $X \subset \mathbb{C}^n$ by the Hironaka algorithm, the degree d , the number l of the polynomials occurring, the embedding dimension n , the number r of the blow-ups and the number q of the affine neighborhoods satisfy, for fixed $m = \dim X$, the recursive equalities above, and are majorized by a function*

$$(r, m, d, n, l, q, \mu) := \Gamma^{(m)}(0, m, d_0, n_0, l_0, q_0, \mu) \in \mathcal{E}^{m+3},$$

for the initial values $d = d_0, n = n_0, l = l_0, q = q_0$.

6.4. Complexity of the algorithm. The principal complexity result of the paper is the following assertion.

THEOREM 6.4.1. *When resolving a marked ideal (X, \mathcal{I}, E, μ) on $X \subset \mathbb{C}^n$ by the Hironaka algorithm, its complexity can be bounded, for fixed $m = \dim X$, by*

$$b^{O(1)} \cdot \mathcal{F}^{(m)}(d_0, n_0, l_0, q_0, \mu),$$

for a certain function $\mathcal{F}^{(m)}(d_0, n_0, l_0, q_0, \mu) \in \mathcal{E}^{m+3}$.

Proof. Indeed, each step of the algorithm consists of solving a certain subroutine (basically, solving a linear system) over the coefficients of the current polynomials. Therefore, Corollary 6.3.1 provides a bound on the number of arithmetic operations with the coefficients of the current polynomials (providing a function from class \mathcal{E}^{m+3}). On the other hand, all the coefficients of the polynomials for the next step are obtained as results of these arithmetic operations, so the bit sizes of the coefficients and the complexity grow by at most $b \cdot \mathcal{F}^{(m)}$, for a function $\mathcal{F}^{(m)} = \mathcal{G}(\Gamma^{(m)}(0, m, d_0, n_0, l_0, q_0, \mu)) \in \mathcal{E}^{m+3}$ with a suitable function $\mathcal{G} \in \mathcal{E}^3$. The cost of a single arithmetic operation is obviously polynomial.

As a corollary we obtain the following theorem.

THEOREM 6.4.2. *When either*

*1. resolving singularities of $X \subset \mathbb{C}^{n_0}$,
or (2) principalizing an ideal sheaf \mathcal{I} on a nonsingular $X \subset \mathbb{C}^{n_0}$,
by the Hironaka algorithm, the degree d , the number l of the polynomials occurring, the embedding dimension n , the number r of the blow-ups, and the number q of the affine neighborhoods satisfy, for fixed $m = \dim X$, the recursive equalities above and are majorized by a function*

$$(r, m, d, n, l, q, 1) := \Gamma^{(m)}(0, m, d_0, n_0, l_0, q_0, 1) \in \mathcal{E}^{m+3},$$

for the initial values $d = d_0, n = n_0, l = l_0, q = q_0$.

The complexity of the algorithm is bounded by

$$b^{O(1)} \cdot \mathcal{F}^{(m)}(d_0, n_0, l_0, q_0, \mu),$$

for a certain function $\mathcal{F}^{(m)}(d_0, n_0, l_0, q_0, \mu) \in \mathcal{E}^{m+3}$.

REMARK. In the proof above, we gave a more explicit form of $\mathcal{F}^{(m)}$ providing additional information on its dependence on r, d, n, l, q, μ . But the main consequence of the Theorem is that $m = \dim X$ provides the most significant contribution to the complexity bound.

REFERENCES

- [1] S. S. ABHYANKAR, *Desingularization of plane curves*, in Algebraic Geometry, Arcata 1981, Proc. Symp. Pure Appl. Math. 40. Amer. Math. Soc., 1983.
- [2] S. S. ABHYANKAR, *Good points of a hypersurface*, Adv. in Math., 68 (1988), pp. 87–256.
- [3] D. ABRAMOVICH AND A. J. DE JONG, *Smoothness, semistability, and toroidal geometry*, J. Alg. Geom., 6 (1997), pp. 789–801.
- [4] D. ABRAMOVICH AND J. WANG, *Equivariant resolution of singularities in characteristic 0*, Math. Res. Letters, 4 (1997), pp. 427–433.
- [5] J. M. AROCA, H. HIRONAKA, AND J. L. VICENTE, *Theory of maximal contact*, Memo Math. del Inst. Jorge Juan, 29, 1975.
- [6] J. M. AROCA, H. HIRONAKA, AND J. L. VICENTE, *Desingularization theorems*, Memo Math. del Inst. Jorge Juan, 30, 1977.
- [7] E. BIERSTONE AND P. MILMAN, *Semianalytic and subanalytic sets*, Publ. Math. IHES, 67 (1988), pp. 5–42.
- [8] E. BIERSTONE AND P. MILMAN, *Uniformization of analytic spaces*, J. Amer. Math. Soc., 2 (1989), pp. 801–836.
- [9] E. BIERSTONE AND P. MILMAN, *A simple constructive proof of canonical resolution of singularities*, in T. Mora and C. Traverso, editors, Effective methods in algebraic geometry, pp. 11–30. Birkhuser, 1991.
- [10] E. BIERSTONE AND D. MILMAN, *Canonical desingularization in characteristic zero by blowing up the maximum strata of a local invariant*, Invent. math., 128 (1997), pp. 207–302.
- [11] E. BIERSTONE AND D. MILMAN, *Desingularization algorithms I. Role of exceptional divisors*, Moscow Math. J., 3 (2003), pp. 751–805.
- [12] E. BIERSTONE AND D. MILMAN, *Functoriality in resolution of singularities*, Publ. RIMS, 44:2 (2008), pp. 609–639.
- [13] R. BLANCO, *Complexity of Villamayor’s algorithm in the non-exceptional monomial case*, Internat. J. Math., 20:6 (2009), pp. 659–678.
- [14] A. BRAVO AND O. VILLAMAYOR, *A strengthening of resolution of singularities in characteristic zero*, Proc. London Math. Soc. (3), 86 (2003), pp. 327–357.
- [15] F. BOGOMOLOV AND T. PANTEV, *Weak Hironaka theorem*, Math. Res. Letters, 3:3 (1996), pp. 299–309.
- [16] D. BROWNAWELL, *Bounds for the degrees in the Nullstellensatz*, Ann. Math. (2), 126 (1987), pp. 577–591.
- [17] V. COSSART, *Desingularization of embedded excellent surfaces*, Thoku Math. J., 33 (1981), pp. 25–33.
- [18] S. D. CUTKOSKY, *Resolution of singularities*, Graduate studies in mathematics, ISSN 1065-7339: v. 63
- [19] S. ENCINAS AND H. HAUSER, *Strong resolution of singularities in characteristic zero*, Comment. Math. Helv., 77 (2002), pp. 821–845.
- [20] S. ENCINAS AND O. VILLAMAYOR, *Good points and constructive resolution of singularities*, Acta Math., 181 (1998), pp. 109–158.
- [21] S. ENCINAS AND O. VILLAMAYOR, *A course on constructive desingularization and equivariance*, in H. Hauser, J. Lipman, F. Oort, and A. Quir, editors, Resolution of Singularities, A research textbook in tribute to Oscar Zariski, volume 181 of Progress in Mathematics. Birkhuser, 2000.
- [22] S. ENCINAS AND O. VILLAMAYOR, *A new theorem of desingularization over fields of characteristic zero*, preprint, 2001.
- [23] J. GIRAUD, *Sur la théorie du contact maximal*, Math. Zeit., 137 (1974), pp. 285–310.
- [24] M. GIUSTI, J. HEINTZ AND J. SABIA, *On the efficiency of effective Nullstellensätze*, Comput. Complexity, 3 (1993), pp. 56–95.
- [25] M. GIUSTI, *Some effectivity problems in polynomial ideal theory*, EUROSAM 84 (Cambridge, 1984), pp. 159–171, Lecture Notes in Comput. Sci., 174, Springer, Berlin, 1984.

- [26] R. GOLDIN AND B. TEISSIER, *Resolving singularities of plane analytic branches with one toric morphism*, preprint ENS Paris, 1995.
- [27] D. GRIGORIEV, *Computational complexity in polynomial algebra*, Proc. Intern. Congress Math., Berkeley, 2 (1986), pp. 1452–1460.
- [28] A. GRZEGORCZYK, *Some classes of recursive functions*, Rozprawy Matematyczne, 4 (1953), pp. 1–44.
- [29] R. HARTSHORNE, *Algebraic Geometry*, Graduate Texts in Mathematics 52, Springer-Verlag, 1977.
- [30] H. HAUSER, *The Hironaka theorem on resolution of singularities (or: A proof we always wanted to understand)*, Bull. Amer. Math. Soc., 40 (2003), pp. 323–403.
- [31] H. HIRONAKA, *An example of a non-Kählerian complex-analytic deformation of Kählerian complex structure*, Annals of Math. (2), 75 (1962), pp. 190–208.
- [32] H. HIRONAKA, *Resolution of singularities of an algebraic variety over a field of characteristic zero*, Annals of Math., 79 (1964), pp. 109–326.
- [33] H. HIRONAKA, *Introduction to the theory of infinitely near singular points*, Memo Math. del Inst. Jorge Juan, 28 (1974).
- [34] H. HIRONAKA, *Idealistic exponents of singularity*, in Algebraic Geometry. The Johns Hopkins centennial lectures, pp. 52–125, Johns Hopkins University Press, Baltimore, 1977.
- [35] Z. JELONEK, *On the effective Nullstellensatz*, Invent. Math., 162 (2005), pp. 1–17.
- [36] A. J. DE JONG, *Smoothness, semistability, and alterations*, Publ. Math. I.H.E.S., 83 (1996), pp. 51–93.
- [37] J. KOLLÁR, *Sharp effective Nullstellensatz*, J. Amer. Math. Soc., 1 (1988), pp. 963–975.
- [38] J. KOLLÁR, *Lectures on Resolution of singularities*, Princeton University Press, 2007.
- [39] D. KOZEN, *Efficient resolution of singularities of plane curves*, Lecture Notes in Comput. Sci., 880 (1994), pp. 1–11.
- [40] J. LIPMAN, *Introduction to the resolution of singularities*, in Arcata 1974, volume 29 of Proc. Symp. Pure Math, pp. 187–229, 1975.
- [41] K. MATSUKI, *Notes on the inductive algorithm of resolution of singularities*, preprint.
- [42] M. MÖLLER AND T. MORA, *Upper and lower bounds for the degree of Groebner bases*, Lect. Notes Comput. Sci., 174 (1984), pp. 172–183.
- [43] T. ODA, *Infinitely very near singular points*, Adv. Studies Pure Math., 8 (1986), pp. :363–404.
- [44] A. SEIDENBERG, *On the length of a Hilbert ascending chain*, Proc. Amer. Math. Soc., 29 (1971), pp. 443–450.
- [45] A. SEIDENBERG, *Constructions in algebra*, Trans. Amer. Math. Soc., 197 (1974), pp. 273–313.
- [46] S. SIMPSON, *Ordinal numbers and the Hilbert Basis Theorem*, J. Symb. Logic, 53 (1988), pp. 961–974.
- [47] J. TEITELBAUM, *On the computational complexity of the resolution of plane curve singularities*, Lecture Notes in Comput. Sci., 358 (1989) Springer, 285–292.
- [48] M. TSFASMAN AND S. VLADUTS, *Algebraic-geometric codes. Mathematics and its Applications*, 58, Kluwer, 1991.
- [49] O. VILLAMAYOR, *Constructiveness of Hironaka’s resolution*, Ann. Scient. Ecole Norm. Sup., 4 (1989), pp. 22:1–32.
- [50] O. VILLAMAYOR, *Patching local uniformizations*, Ann. Scient. Ecole Norm. Sup. 4, 25 (1992), pp. 629–677.
- [51] O. VILLAMAYOR, *Introduction to the algorithm of resolution*, in Algebraic geometry and singularities, La Rabida 1991, pages 123–154. Birkhuser, 1996.
- [52] K. WAGNER AND G. WECHSUNG, *Computational Complexity*, Mathematics and its Applications, 21, 1986.
- [53] J. WŁODARCZYK, *Simple Hironaka resolution in characteristic zero*, J. Amer. Math. Soc., 18 (2005), pp. 779–822